

Heegner points on Elliptic curves

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Diophantine equations

Central problem in Number Theory

- Polynomial equations with rational coefficients
- Interested in: rational solutions
- $C: x^2 + y^2 - 1 = 0$
 - ▶ $C(\mathbb{Q}) = \{ \text{solutions of } C \text{ with rational coordinates} \} \subset \mathbb{A}^2$
 - ▶ $(1, 0) \in C(\mathbb{Q})$
 - ▶ In fact, this equation has infinitely many solutions

Open question

Is there an algorithm that given a diophantine equation C computes $C(\mathbb{Q})$?

- Two variables: $f(x, y) = 0$ with $f \in \mathbb{Q}[x, y]$ is called a plane curve
 - ▶ This case is already open
 - ▶ But a lot is known: for example, topology plays a role!

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- $C: f(x, y) = 0 \rightsquigarrow C(\mathbb{C}) = \{ \text{solutions with } x, y \in \mathbb{C} \}$
 - ▶ Write $x = a + bi$, $y = c + di$ with $a, b, c, d \in \mathbb{R}$
 - ▶ 2 equations in 4 unknowns \rightsquigarrow it is a surface (assume nonsingular)
 - ▶ non-compact, but can be compactified adding finitely many points:

$$\bar{C}: z^d f\left(\frac{x}{z}, \frac{y}{z}\right) = 0 \text{ and } \bar{C}(\mathbb{Q}) \subset \mathbb{P}^2$$

- $\bar{C}(\mathbb{C})$ is homeomorphic to a g -holed torus, where g is the genus
- Formula: $g = (d - 1)(d - 2)/2$ where d is the degree of f .
- $C: x^2 + y^2 - 1 = 0 \rightsquigarrow \bar{C}(\mathbb{C})$ is homeomorphic to a sphere

Falting's Theorem, 1984 (a.k.a. Mordell Conjecture)

If $g > 1$ then $C(\mathbb{Q})$ is finite.

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Elliptic curves

Elliptic curve

A non-singular genus 1 curve which has a rational point.

- By a rational change of variables they can be transformed into

$$E: y^2 = x^3 + Ax + B, \text{ with } A, B \in \mathbb{Z}, A^3 - 27B^2 \neq 0.$$

- $[0: 1: 0]$ point of $y^2z = x^3 + Axz^2 + Bz^3$, the only point at infinity
- Example:

$$E: y^2 = x^3 - 16x + 16$$

$$\triangleright (0, 4), (4, 4) \in E(\mathbb{Q})$$

Key property

One can define a group operation

$$\begin{aligned} +: E(\mathbb{Q}) \times E(\mathbb{Q}) &\longrightarrow E(\mathbb{Q}) \\ (P, Q) &\longmapsto P + Q \end{aligned}$$

- From $P = (0, 4)$ and $Q = (4, 4) \rightsquigarrow P + Q \in E(\mathbb{Q})$

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- $[0: 1: 0]$ point of $y^2z = x^3 + Axz^2 + Bz^3$, the only point at infinity
- Example:

$$E: y^2 = x^3 - 16x + 16$$

▶ $(0, 4), (4, 4) \in E(\mathbb{Q})$

Key property

One can define a group operation

$$\begin{aligned} +: E(\mathbb{Q}) \times E(\mathbb{Q}) &\longrightarrow E(\mathbb{Q}) \\ (P, Q) &\longmapsto P + Q \end{aligned}$$

- From $P = (0, 4)$ and $Q = (4, 4) \rightsquigarrow P + Q \in E(\mathbb{Q})$

Elliptic curves

Elliptic curve

A non-singular genus 1 curve which has a rational point.

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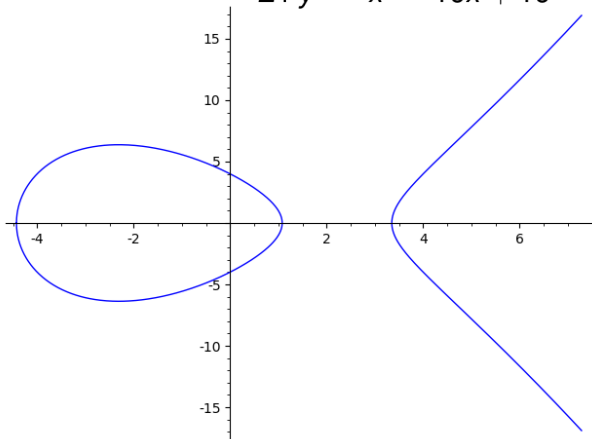
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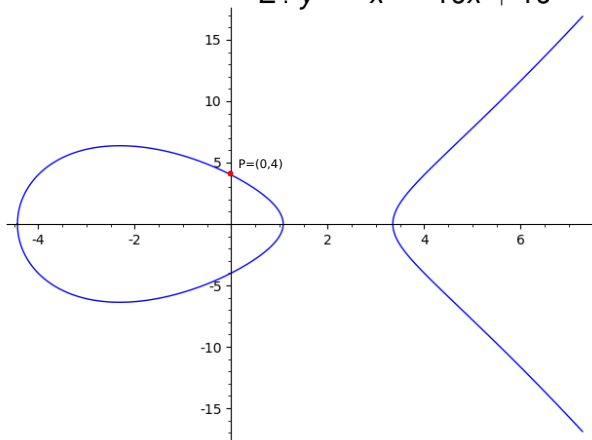
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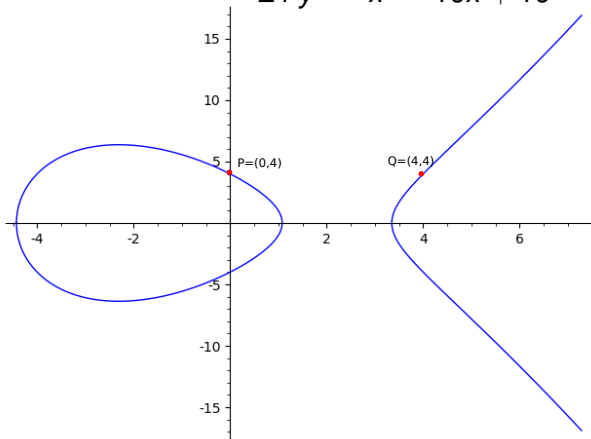
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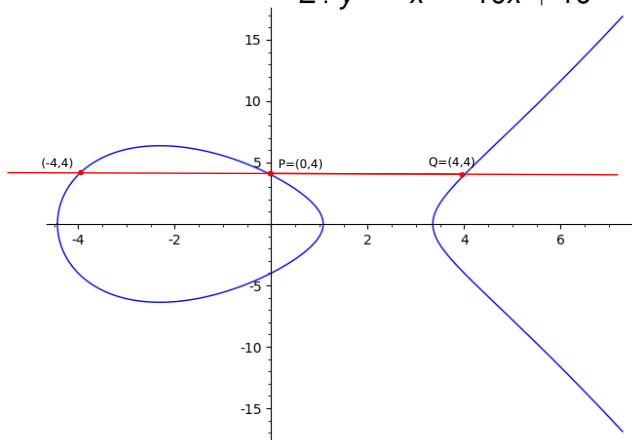
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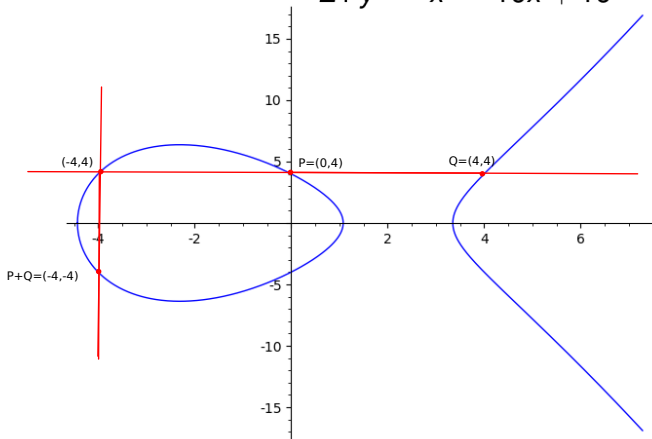
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The group of points

- This makes $E(\mathbb{Q})$ into an abelian group.

Mordell Theorem (1922)

$E(\mathbb{Q})$ is a finitely generated abelian group.

- Structure Theorem: $E(\mathbb{Q}) \simeq \mathbb{Z}^r \oplus T$
 - ▶ T is finite, its points have finite order
 - ▶ r is called the rank of E .
- T is pretty well understood:

Theorem (Mazur, 1977)

$$T \simeq \begin{cases} \mathbb{Z}/N\mathbb{Z}, & 1 \leq N \leq 10 \text{ or } N = 12 \\ \mathbb{Z}/2\mathbb{Z} \times \mathbb{Z}/2N\mathbb{Z}, & 1 \leq N \leq 4. \end{cases}$$

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- $E: y^2 = x^3 + Ax + B$ with $A, B \in \mathbb{Z}$
- For every prime number p :

$$y^2 \equiv x^3 + Ax + B \pmod{p}$$

- Define $a_p := p - \#\text{solutions}$
- $L(E, s) = \prod_{p|N} \frac{1}{1 - a_p p^{-s} + p^{1-2s}} = \sum_{n \geq 1} a_n n^{-s}$
 - ▶ Product only converges for $\text{Re}(s) > 3/2$, but can be extended to \mathbb{C}

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The rank r of E equals $\text{ord}_{s=1} L(E, s)$

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If $\text{ord}_{s=1} L(E, s)$ is 0 or 1, the BSD Conjecture is true.

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Complex analysis meets number theory

- $E: y^2 = x^3 + Ax + B \rightsquigarrow E(\mathbb{C})$ is homeomorphic to a torus

Weierstrass Uniformization Theorem

There is a lattice $\Lambda_E \subset \mathbb{C}$ such that $E(\mathbb{C}) \simeq \mathbb{C}/\Lambda_E$ as complex varieties.

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- $L(E, s) = \sum a_n n^{-s}$ with $a_p := p - \#\text{solutions (mod } p)$
- $f_E(z) := \sum_{n \geq 1} a_n e^{2\pi i n z}$ converges for $\text{Im}(z) > 0$

Modularity Theorem (Wiles, Breuil-Conrad-Diamond-Taylor)

$f_E(z)$ is a modular form.

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- Associated to imaginary quadratic numbers $w = a + b\sqrt{-D}$
- $P_w = \int_{i\infty}^w f_E(z) dz \in \mathbb{C}/\Lambda_E \simeq E(\mathbb{C})$

Complex multiplication + Gross-Zagier

If $\text{ord}_{s=1} L(E, s) = 1$ choosing w appropriately $P_w \in E(\mathbb{Q})$ and is of infinite order.

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Stark-Heegner points

Natural question

Are there points P_w associated to **real** quadratic numbers $a + b\sqrt{D}$?

- Henri Darmon in 2000 proposed a construction using *p*-adic integrals instead of complex integrals.
- These are called Stark–Heegner (or Darmon) points, and have been constructed also for curves over other number fields.
- Rationality of the points is still conjectural, but the method can be used to compute points in practice: very efficient algorithm

Example (G.-Masdeu)

- $E: y^2 + xy = x^3 - 8x$
- $\left(\frac{3449809443179}{499880696975}, -\frac{3449809443179}{999761793950} + \frac{3600393040902501011}{3935597293546963250} \sqrt{341} \right) \in E(\mathbb{Q}(\sqrt{341}))$

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Stark-Heegner points

Natural question

Are there points P_w associated to **real** quadratic numbers $a + b\sqrt{D}$?

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Heegner points on Elliptic curves

Xevi Guitart (UB)

BMS-BGSMath Junior Meeting 2022