

Effective Algorithms in Harmonic Analysis and Applications to Signal Processing

Master in Advanced and Professional Mathematics Course 2010/2011

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Introduction

The interest in Fourier and Wavelets Transforms is because Harmonic Analysis is very useful in signal processing nowadays. Signal compression, de-noising, recognition are some examples of its range. In Mathematics, we know its development in continuous (\mathbb{R}) and discrete time (\mathbb{Z}), but here we will deal only with the discrete part. All the signals will be treated like vectors and matrices. We will not only deal with Harmonic Analysis, but also with other fields in signal processing like Morphological operations and Image filtering. All these topics have a theoretical part based on rigorous mathematical theories, that will be explained in detail. But our goal is to go further by implementing, on a computer, the most important techniques and algorithms used to develop in practice these methods.

In first and second chapter, Discrete Fourier Transform and Fast Fourier Transform respectively, we will prove why is used the fast algorithm to compute the Fourier transform. The fact is that, \bigcirc compute the discrete transform, the time required is of the order of N^2 , where N is the length of the signal; but if we pute the fast algorithm, the time spend in it will be of $N \log N$. For example, if we spent \bigcirc to compute the Discrete Fourier Transform, we will spend \bigcirc in \bigcirc Fast Fourier Transform.



Morphological operations are used to determine features of the analysed objects, their for heir neighborhood. In the chapter of Intensity Images and Filtering we will work in the space domain and in the frequency domain, because in this moment we will have the necessary tools to compute the Fourier transform.

In each chapter there is a section of *Matlab Implementation*. All the functions



and GUI's (Graphical user interface's) have been implemented with the program Matlab (Matrix Laboratory). The *IDE* (Integrated development environment) used in the project has been *Microsoft Visual C++ 2008 Express Edition* and *Matlab R2009b*. The languages supported are C/C++ and *Matlab*, its own language, respectively.

Matlab has a compile library in C++ to compute the Fast Fourier Transform, so for this reason, we use Microsoft Visual. First, the code is implemented in C++with the libraries OpenCV (Open Computer Vision) and the header of Matlab, mex.h. Then a mex file is created from Matlab to make the code works on it. This new file works like another function in the program. The code of each function is at the end of its corresponding chapter to clarify its development. The GUI's code has not been included since it is to large, around 200 lines of code for each one.

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Chapter 1

The Discrete Fourier Transform

In this chapter we will introduce and deal with the discrete Fourier transform. First, we will introduce the notation, definitions and properties, followed by its matrix representation, its code implementation in one and two dimensions, and finally some examples of the implemented algorithms.

1.1 Properties of the Discrete Fourier Transform

Notation 1.1. In the discrete case we work with \mathbb{C}^N vectors, sequences of N complex numbers. See [1].

- 1. Now we consider z as a function defined in the finite set $\mathbb{Z}_N = \{0, ..., N-1\}$. We will write the z's components in this way z = (z(0), ..., z(N-1)).
- 2. For the finite dimensional case we introduce the vector space over \mathbb{C}

$$\ell^2(\mathbb{Z}_N) = \{ z \mid z(j) \in \mathbb{C}, \ 0 \le j \le N-1 \}.$$

Note that the Euclidean basis is an orthogonal basis of this space.

3. Inner product

$$\langle z, w \rangle = \sum_{k=0}^{N-1} z(k) \overline{w(k)}$$

4. $\ell^2(\mathbb{Z}_N) - norm$

$$||z|| = \left(\sum_{k=0}^{N-1} |z(k)|^2\right)^{1/2}$$

5. We can extend the sequence of z to any $j \in \mathbb{Z}$ assuming that z is periodic with period N:

$$z(j+N) = z(j), \text{ for all } j \in \mathbb{Z}.$$

The definition and lemma below are necessary for introducing the Discrete Fourier Transform (DFT).

Definition 1.2. Define $E_0, ..., E_{N-1} \in \ell^2(\mathbb{Z}_N)$ by

(1.1)
$$E_m(n) = \frac{1}{\sqrt{N}} e^{2\pi i m n/N}, \ 0 \le m, n \le N - 1.$$

Lemma 1.3. The set $\{E_0, ..., E_{N-1}\}$ is an orthonormal basis of $\ell^2(\mathbb{Z}_N)$.

Proof. Suppose $m, m' \in \{0, ..., N-1\}$. Then

$$\langle E_m, E_{m'} \rangle = \sum_{n=0}^{N-1} E_m(n) \overline{E_{m'}(n)} = \sum_{n=0}^{N-1} \frac{1}{\sqrt{N}} e^{2\pi i m n/N} \frac{1}{\sqrt{N}} e^{-2\pi i m' n/N}$$
$$= \frac{1}{N} \sum_{n=0}^{N-1} e^{2\pi i (m-m')n/N}.$$

If m = m'

$$\langle E_m, E_m \rangle = \frac{1}{N} \sum_{n=0}^{N-1} e^{2\pi i (m-m)n/N} = \frac{1}{N} \sum_{n=0}^{N-1} 1 = 1.$$

If $m \neq m'$

$$\langle E_m, E_{m'} \rangle = \frac{1}{N} \sum_{n=0}^{N-1} e^{2\pi i (m-m')n/N} = \frac{1}{N} \sum_{n=0}^{N-1} \left(e^{2\pi i (m-m')/N} \right)^n$$

= $\frac{1}{N} \frac{\left(e^{2\pi i (m-m')/N} \right)^{N-1} e^{2\pi i (m-m')/N} - 1}{e^{2\pi i (m-m')/N} - 1} = \frac{1}{N} \frac{e^{2\pi i (m-m')} - 1}{e^{2\pi i (m-m')/N} - 1} = 0.$

Since $\{E_0, ..., E_{N-1}\}$ is an orthonormal basis for $\ell^2(\mathbb{Z}_N)$, we have for all $z, w \in \ell^2(\mathbb{Z}_N)$ the following properties:

(1.2)
$$z = \sum_{m=0}^{N-1} \langle z, E_m \rangle E_m$$

(1.3)
$$\langle z, w \rangle = \sum_{m=0}^{N-1} \langle z, E_m \rangle \,\overline{\langle w, E_m \rangle}$$

(1.4)
$$||z||^2 = \sum_{m=0}^{N-1} |\langle z, E_m \rangle|^2$$

where

(1.5)
$$\langle z, E_m \rangle = \sum_{m=0}^{N-1} z(n) \overline{\frac{1}{\sqrt{N}} e^{2\pi i m n/N}} = \frac{1}{\sqrt{N}} \sum_{m=0}^{N-1} z(n) e^{-2\pi i m n/N}$$

Definition 1.4. Suppose $z = (z(0), ..., z(N-1)) \in \ell^2(\mathbb{Z}_N)$. For m = 0, ..., N-1 define

(1.6)
$$\hat{z}(m) = \sum_{n=0}^{N-1} z(n) e^{-2\pi i m n/N}.$$

Let

(1.7)
$$\hat{z} = (\hat{z}(0), ..., \hat{z}(N-1)).$$

Then $\hat{z} \in \ell^2(\mathbb{Z}_N)$ and the map $z \longrightarrow \hat{z}$ is called the discrete Fourier transform *DFT*.

Remark 1.5. It is easy to see that the map

$$\ell^2 \left(\mathbb{Z}_N \right) \stackrel{\sim}{\longrightarrow} \ell^2 \left(\mathbb{Z}_N \right)$$
$$z \longmapsto \hat{z}$$

is a linear transformation.

We observe that \hat{z} is also an N-periodic vector:

$$\hat{z}(m+N) = \sum_{n=0}^{N-1} z(n) e^{-2\pi i (m+N)n/N} = \sum_{n=0}^{N-1} z(n) e^{-2\pi i mn/N} e^{-2\pi i Nn/N} = \hat{z}(m).$$

Note that

(1.8)
$$\hat{z}(m) = \sqrt{N} \langle z, E_m \rangle.$$

Theorem 1.6. Let $z = (z(0), ..., z(N-1)), w = (w(0), ..., w(N-1)) \in \ell^2(\mathbb{Z}_N).$ Then

1. Fourier inversion formula:

(1.9)
$$z(n) = \frac{1}{N} \sum_{m=0}^{N-1} \hat{z}(m) e^{2\pi i m n/N}, \text{ for } n = 0, ..., N-1.$$

2. Parseval's relation:

(1.10)
$$\langle z, w \rangle = \frac{1}{N} \sum_{m=0}^{N-1} \hat{z} \overline{\hat{w}} = \frac{1}{N} \langle \hat{z}, \hat{w} \rangle.$$

3. Plancherel's formula:

(1.11)
$$||z||^2 = \frac{1}{N} \sum_{m=0}^{N-1} |\hat{z}(m)|^2 = \frac{1}{N} ||\hat{z}||^2.$$

Proof.

1.

$$z(n) = \sum_{m=0}^{N-1} \langle z, E_m \rangle E_m(n) = \sum_{m=0}^{N-1} \frac{1}{\sqrt{N}} \hat{z}(m) \frac{1}{\sqrt{N}} e^{2\pi i m n/N} = \frac{1}{N} \sum_{m=0}^{N-1} \hat{z}(m) e^{2\pi i m n/N}.$$

$$\langle z, w \rangle = \sum_{m=0}^{N-1} \langle z, E_m \rangle \,\overline{\langle w, E_m \rangle} = \sum_{m=0}^{N-1} \frac{1}{\sqrt{N}} \hat{z}\left(m\right) \overline{\frac{1}{\sqrt{N}} \hat{w}\left(m\right)} = \frac{1}{N} \sum_{m=0}^{N-1} \hat{z}\left(m\right) \overline{\hat{w}\left(m\right)}.$$

3.

$$||z||^{2} = \sum_{m=0}^{N-1} |\langle z, E_{m} \rangle|^{2} = \frac{1}{N} \sum_{m=0}^{N-1} |\hat{z}(m)|^{2}.$$

Remark 1.7.

- 1. We call $\{E_0, ..., E_{N-1}\}$ the Fourier basis for $\ell^2(\mathbb{Z}_N)$.
- 2. The vector representing z with respect to the Fourier basis is \hat{z} .
- 3. The Fourier inversion formula in the above theorem is the change of basis formula for the Fourier basis.
- 4. The DFT components $\hat{z}(m)$ are the components of z in the Fourier basis.

Definition 1.8. The convolution of $z, w \in \ell^2(\mathbb{Z}_N)$ is defined as

(1.12)
$$(z*w)(n) = \sum_{l=0}^{N-1} z(l) \cdot w(n-l), \quad n = 0, ..., N-1.$$

Theorem 1.9. Suppose $z, w \in \ell^2(\mathbb{Z}_N)$. Then,

(1.13)
$$(z * w)^{\wedge}(n) = \hat{z}(n) \cdot \hat{w}(n), \quad n = 0, ..., N - 1.$$

Proof.

$$(z * w)^{\wedge}(n) = \sum_{m=0}^{N-1} \left(\sum_{l=0}^{N-1} z(l) \cdot w(m-l) \right) e^{-2\pi i m n/N}$$

= $\sum_{l=0}^{N-1} z(l) \left(\sum_{m=0}^{N-1} w(m-l) \cdot e^{-2\pi i (m-l)n/N} \right) e^{-2\pi i l n/N}$
= $\left(\sum_{l=0}^{N-1} z(l) \cdot e^{-2\pi i l n/N} \right) \left(\sum_{m=0}^{N-1} w(m) \cdot e^{-2\pi i m n/N} \right)$
= $\hat{z}(n) \cdot \hat{w}(n).$

1.1.1 Matrix representation of DFT

Notation 1.10.

- $\omega_N = e^{-2\pi i/N}$.

$$-\omega_N^{mn} = e^{-2\pi i m n/N}.$$

-
$$\hat{z}(m) = \sum_{m=0}^{N-1} z(n) \,\omega_N^{mn}$$

Definition 1.11. Let W_N be the matrix $(w_{mn})_{0 \le m, n \le N-1}$ such that $w_{mn} = \omega_N^{mn}$

(′ 1	1		1	١
	1	ω_N		ω_N^{N-1}	
	÷	÷	۰.	÷	
	1	ω_N^{N-1}		$\omega_N^{(N-1)(N-1)}$	

Regarding $z, \hat{z} \in \ell^2(\mathbb{Z}_N)$ as column vectors, we can write

$$\hat{z} = W_N \cdot z$$

Signal processing deals with images, sounds and others, so we will refer to images and sounds like signals.

Definition 1.12. Let us consider a signal f of size $M \times N$. We define the 2D Discrete Fourier Transform

(1.14)
$$\hat{f}(u,v) = \sum_{j=0}^{M-1} \sum_{k=0}^{N-1} f(j,k) e^{-2\pi i (uj/M + vk/N)},$$

where u = 0, ..., M - 1 and v = 0, ..., N - 1.

1.2 DFT Matlab implementation

Implementation of DFT using the W's matrix.

```
_ Code _
function Z = DFT_mat(z)
N = length(z);
W = ones(N);
w = \exp(-2*pi*i/N);
% Auxiliar array&matrix %
x = 1: N-1;
mat = zeros(N-1);
for k=1:N-1
   mat(k,:) = x.*k;
end
% %% Matrix of w's %%
W(2:end, 2:end) = w.^{(mat)};
%% Check if z is a column vector %%
[r,c] = size(z);
if c == N
        Z = W * z';
else
        Z = W * z;
end
```

Implementation of DFT using the definition of the map

$$\hat{z}(m) = \sum_{n=0}^{N-1} z(n) e^{-2\pi i m n/N}$$

function Z=DFT(z)

N = length(z);

Z = zeros(1,N);

for m=1:N

for n=1:N

Z(m) = Z(m) + z(n)*exp(-2*pi*i*((m-1)*(n-1)/N));

end

end

2D DFT Matlab implementation.

```
Code
function F=DFT2(signal)
[r,c] = size(signal);
F = zeros(r,c);
for u=1:r
for v=1:c
for j=1:r
for k=1:c
F(u,v) = F(u,v) + signal(j,k)*exp(-2*pi*i*((v-1)*(k-1)/c)
+ (j-1)*(u-1)/r);
end
end
```

end

 $\quad \text{end} \quad$

1.3 Examples

Definition 1.13. The spectrum of the Fourier Transform of a signal is its absolute values.

Let us consider,



Figure 1.2: DFT of $z_1(n) = \cos(2\pi 15n/16)$



Figure 1.3: DFT of $z_2(n) = \sin(\pi n^2/256)$

4. $z_3(n) = 3\sin(2\pi7n/512) - 4\cos(2\pi8n/512) \in \ell^2(\mathbb{Z}_{512}).$



Figure 1.4: DFT of $z_3(n) = 3\sin(2\pi 7n/512) - 4\cos(2\pi 8n/512)$

The following table represents the time (seconds) spent in processing the signals above. The first column DFTmat and the second column DFT represent the algorithms 1.1.1 and 1.6 respectively.

Length	DFTmat (s)	DFT (s)
4	$1.5142 \cdot 10^{-4}$	$9.1352 \cdot 10^{-5}$
16	$6.8640 \cdot 10^{-4}$	$4.1234 \cdot 10^{-4}$
128	0.0654	0.0179
512	0.6629	0.1993
1024	2.8469	0.6662

Figure 1.5 represents the graph of time varying the length of the signals. It can be observe that the algorithm that uses the matrix to calculate the Fourier Transform is slower than the one that applies the definition directly.



Figure 1.5: Time comparison

Chapter 2

The Fast Fourier Transform

In this chapter we will introduce the Fast Fourier Transform to compute Fourier definition explained above, to improve the execution time in computers.

2.1 Description of FFT

The *DFT* can be computed by a fast algorithm, known as the *fast Fourier trans*form (*FFT*). We will see the efficiency of *FFT* by comparing the number of complex multiplications in each algorithm, *DFT* vs *FFT*. We do not make reference to the number of complex additions because complex multiplications have slowly time computation. We have $\hat{z} = W_N \cdot z$, where the matrix W_N has dimensions $N \times N$, and hence the number of complex multiplications is N^2 . See [3].

We consider only the case in which the length N of the vector is even, and in particular a power of 2. The lemma below shows the basic idea behind the FFT.

Lemma 2.1. Suppose $M \in \mathbb{N}$, and N = 2M. Let $z \in \ell^2(\mathbb{Z}_N)$. Define $u, v \in \ell^2(\mathbb{Z}_M)$ by

$$u(k) = z(2k)$$
 for $k = 0, ..., M - 1$,
 $v(k) = z(2k + 1)$ for $k = 0, ..., M - 1$.

In other words,

$$u = (z(0), z(2), ..., z(N-4), z(N-2)),$$
$$v = (z(1), z(3), ..., z(N-3), z(N-1)).$$

Let \hat{z} denote the DFT of z, given by $\hat{z} = W_N z$. Let \hat{u} and \hat{v} denote the DFTs of u and v respectively, defined for M = N/2, $\hat{u} = W_M \cdot u$, $\hat{v} = W_M \cdot v$. Then, for m = 0, ..., M - 1,

(2.1)
$$\hat{z}(m) = \hat{u}(m) + e^{-2\pi i m/N} \hat{v}(m);$$

and for m = M, ..., N - 1 and l = m - M, we have

(2.2)
$$\hat{z}(l+M) = \hat{u}(l) - e^{-2\pi i l/N} \hat{v}(l).$$

Proof. For any m = 0, ..., N - 1

$$\hat{z}(m) = \sum_{n=0}^{N-1} z(n) e^{-2\pi i m n/N}.$$

The sum over n = 0, ..., N-1 can be broken into the sum over the even values n = 2k, k = 0, ..., M-1 plus the sum over the odd values n = 2k + 1, k = 0, ..., M-1:

$$\hat{z}(m) = \sum_{k=0}^{M-1} z(2k) e^{-2\pi i 2km/N} + \sum_{k=0}^{M-1} z(2k+1) e^{-2\pi i (2k+1)m/N}$$
$$= \sum_{k=0}^{M-1} u(k) e^{-2\pi i km/(N/2)} + e^{-2\pi i m/N} \sum_{k=0}^{M-1} v(k) e^{-2\pi i km/(N/2)}$$
$$= \sum_{k=0}^{M-1} u(k) e^{-2\pi i km/M} + e^{-2\pi i m/N} \sum_{k=0}^{M-1} v(k) e^{-2\pi i km/M}.$$

In the case of m = 0, ..., M - 1 we obtain

$$\hat{u}(m) + e^{-2\pi i m/N} \hat{v}(m) \,.$$

Suppose m = M, ..., N-1 and l = m - M, so that l = 0, ..., M-1. By writing

m = l + M in the expression 2.3 we get

$$\hat{z}(m) = \sum_{k=0}^{M-1} u(k) e^{-2\pi i k (l+M)/M} + e^{-2\pi i (l+M)/N} \sum_{k=0}^{M-1} v(k) e^{-2\pi i k (l+M)/M}$$
$$= \sum_{k=0}^{M-1} u(k) e^{-2\pi i k l/M} - e^{-2\pi i l/N} \sum_{k=0}^{M-1} v(k) e^{-2\pi i k l/M},$$

since $e^{-2\pi i M/N} = e^{-2\pi i (N/2)/N} = e^{-\pi i} = -1$. Then, in the case m = M, ..., N - 1

$$\hat{u}(m) - e^{-2\pi i m/N} \hat{v}(m)$$

2.2 FFT algorithm: Decimation in Time, Radix2

We will assume that N is a power of 2, $N = 2^n$, with $n \in \mathbb{Z}^+$, Walker [3]. By the Lemma 2.1 we write

$$\hat{z}(m) = \hat{u}(m) + \omega_N^m \cdot \hat{v}(m), \ m = 0, ..., M - 1,$$
$$\hat{z}(m) = \hat{u}(m) - \omega_N^m \cdot \hat{v}(m), \ m = M, ..., N - 1,$$

where $\omega_N^m = e^{-2\pi i m/N}$. These calculations can be plotted as



Figure 2.1: Butterfly diagram.

To apply equations (2.1) and (2.2), we first compute \hat{u} and \hat{v} . Each of these is a vector of length M = N/2, so each can be computed with M^2 complex multiplications.

Then we compute the products $\omega_N^m \cdot \hat{v}(m)$, m = 0, ..., M - 1. This requires M multiplications. So the total number of complex multiplications required to compute \hat{z} is at most

$$2M^{2} + M = 2\left(\frac{N}{2}\right)^{2} + \frac{N}{2} = \frac{1}{2}\left(N^{2} + N\right).$$

For N large, this is essentially $N^2/2$, whereas N^2 multiplications are required to compute \hat{z} directly.

Notation 2.2. We define $\#_N$ as the least number of complex multiplications required to compute the DFT of a vector of length $N \in \mathbb{Z}^+$.

If N = 2M, then equations (2.1) and (2.2) reduce the computation of \hat{z} to the computation of two *DFTs* of size *M*, plus *M* additional complex multiplications

(2.4)
$$\#_N \le 2\#_M + M.$$

Lemma 2.3. Suppose $N = 2^n$ for some $n \in \mathbb{N}$. Then,

$$\#_N \le \frac{1}{2} N \log_2 N$$

Proof. By induction on n:

Let n = 1. In this case we have a vector of length 2^1 , $z = (z_1, z_2)$. Then,

$$\hat{z} = W_2 z = \begin{pmatrix} 1 & 1 \\ 1 & \omega_2 \end{pmatrix} \begin{pmatrix} z_1 \\ z_2 \end{pmatrix} = \begin{pmatrix} z_1 + z_1 \\ z_1 - z_2 \end{pmatrix},$$

where $\omega_2 = e^{-\pi i}$.

This computation does not require any complex multiplication, so

$$\#_2 = 0 < 1 = \frac{1}{2} 2 \log_2 2.$$

By induction, assume it holds for n-1. Then,

$$\#_{2^{n}} \underbrace{\leq}_{(2.4)} 2 \#_{2^{n-1}} + 2^{n-1} \leq 2 \left(\frac{1}{2} 2^{n-1} \log_2 2^{n-1} \right) + 2^{n-1}$$
$$= 2^{n-1} (n-1) + 2^{n-1} = n2^{n-1} = \frac{1}{2} n2^n = \frac{1}{2} n \log_2 n.$$

Example 2.4. Let us consider a vector of length $N = 2^3$.



Figure 2.2: First reduction in FFT algorithm for N = 8.

We can split \hat{u} and \hat{v} into two DFT half-size of the original vector and so on. In this example we will consider the vector z = (1, 2, 0, i, 1, -1, 3, -i). The length of the vector is $8 = 2^3$. This means that we can split \hat{u} and \hat{v} three times to obtain the DFT of vectors of length 1. Figure 2.3 shows the three splits to calculate its FFT.



Figure 2.3: FFT for N = 8.

Example 2.5. Let us consider the same data of the Section 1.3. In the table 2.1 we compare the times spent in the DFT and FFT algorithms.

Length	DFTmat (s)	DFT (s)	fft (s)
4	$6.4813 \cdot 10^{-5}$	$3.9111 \cdot 10^{-5}$	$2.6260 \cdot 10^{-5}$
16	$1.5477 \cdot 10^{-4}$	$1.0392 \cdot 10^{-4}$	$2.7657 \cdot 10^{-5}$
128	0.0103	0.0042	$3.4083 \cdot 10^{-5}$
512	0.1286	0.0648	$4.9168 \cdot 10^{-5}$

Table 2.1: Comparison between DFT and FFT times.

Figure 2.4 represents the time spent in the differents algoritms: DFTmat, DFT and FFT. We represent only the FFT inside a graph because its values cannot be seen in the first graph.



Figure 2.4: Comparison of DFTmat, DFT and FFT algorithms

2.3 Code implementation

2.3.1 FFT Matlab implementation

Auxiliar functions: *padding1D* and *iterative_twiddle*.

```
_____ Code _____
function l_pad = padding1D(n)
1_pad = 2;
while n>l_pad
   l_pad = l_pad*2;
end
             _____ Code _____
function y = iterative_twiddle(x,N,w)
y = bitrevorder(x);
for i=1:log2(N)
   d = pow2(i);
   ktwiddle = 1;
   for k=1:d/2
      for m=k:d:N
          t = w(ktwiddle)*y(m+d/2);
          z = y(m);
          y(m) = z + t;
          y(m+d/2) = z - t;
       end
       ktwiddle = ktwiddle + pow2(log2(N)-i);
```

end

end

Main function $FFT_{-iterative}$

```
_____ Code _____
function FFT = FFT_iterative(signal)
N = length(signal);
%% Check N is a power of 2 %%
if isinteger(log2(N))
   N_pad = N;
else
   N_pad = padding1D(N);
end
w = twiddle(N_pad);
signal_pad = zeros(1,N_pad);
signal_pad(1:N) = signal;
FFT = iterative_twiddle(signal_pad,N_pad,w);
%%
    EXECUTION
\% z = [1 1 1 i 1 -1 1 -i];
%% tic;FFT_iterative(z);t=toc
%% tic;fft(z);t=toc
```

2.3.2 2D FFT Matlab implementation

The main program uses the same auxiliary functions as the one dimensional.

```
Code _
function fft_2 = FFT2_iterative(im)
[rows,cols] = size(im);
%% Check N is a power of 2 %%
if isinteger(log2(rows))
    rows_pad = rows;
else
    rows_pad = padding1D(rows);
end
if isinteger(log2(cols))
    cols_pad = cols;
else
    cols_pad = padding1D(cols);
end
%% Pad image %%
im_pad = zeros(rows_pad,cols_pad);
im_pad(1:rows,1:cols) = im;
wc = twiddle(cols_pad);
wr = twiddle(rows_pad);
 for i=1:rows_pad
     fft(i,:) = iterative_twiddle(im_pad(i,:),cols_pad,wc);
 end
for j=1:cols_pad
    fft_2(:,j) = iterative_twiddle(im_pad(:,j),rows_pad,wr);
```

end

2.3.3 FFT C++ implementation

Twiddles function. In this function we calculate the ω_N values, so we only have to calculate the elements till N/2.

Code
void
FFT::Compute_TWIDDLES(int N)
{
 for(int k = 0; k < N/2; k++)
 twiddles[k] = polar(1.0, -2*pi*(k/(double)N));
}</pre>

Bit reversal function. This function reorders the vector. First, we compute the positions of the vector in bits, and then we swap the bit string. For example, the position 1 in the vector corresponds with the bit string 001, and it becomes the bit string 100. It can be seen better in the Figure 2.3.

```
____Code ______
void
FFT::BitReversal(Mat &mat,int p)
{
    unsigned long x = mat.cols;
    unsigned long h = 0;
    int i = 0;
    int i = 0;
    int k, k_aux;
    complex<double> temp;
```

```
int *v;
v = new int[mat.cols];
Mat m_aux = mat.clone();
for(k = x-1; k >= 0 ;k--){
    k_aux = k;
    for(h = i = 0; i < p; i++){
        h = (h << 1) + (k_aux & 1);
        k_aux >>= 1;
    }
    v[k] = h;
}
for(int i=0; i<mat.cols;i++)
    mat.at<complex<double> >(0,i) =
    m_aux.at<complex<double> >(0,v[i]);
delete[] v;
```

}

FFT function

Code _

}

}

```
// Compute twiddle vector //
Compute_TWIDDLES(cols);
// Log2 of cols //
unsigned int cols_aux = cols;
while ( cols_aux >>= 1) // unroll for more speed...
{
  r++;
}
// bitreversal //
BitReversal(FFT_fft,r);
for(int i=1; i <= r; i++){</pre>
d = 1 << i;
ktwiddle = 0;
 for(int k = 0; k < d/2; k++)
 for(int j = k; j < cols; j+=d){</pre>
  t = (FFT_fft.at<complex<double> >(0,j + d/2)).operator
         *=(twiddles[ktwiddle]);
         z = FFT_fft.at<complex<double> >(0,j);
         FFT_fft.at<complex<double> >(0,j) = z + t;
         FFT_fft.at<complex<double> >(0,j + d/2) = z - t;
 }
         p = 1 << (r-i);
         ktwiddle += p;
 }
```

2.3.4 2D FFT C++ implementation

The bit reversal function and twiddles function used in the code are the same as the implemented in one dimension. See Section 2.3.3.

FFT2D function.

```
Code
void
FFT::Compute_FFT2D()
{
/// FFT1D for all the rows ///
// Compute twiddle vector //
Compute_TWIDDLES(cols);
for(int i=0; i< rows; i++)
Compute_FFT1D_r(FFT_fft.row(i));
/// FFT1D for all the columns ///
// Compute twiddle vector //
Compute_TWIDDLES(rows);
for(int i=0; i< cols; i++)
Compute_FFT1D_c(FFT_fft.col(i));
```

}

{

1D FFT function for rows.

```
_ Code _____
void
        FFT::Compute_FFT1D_r(Mat& fft_r )
        complex<double> z(0.0,0.0), t(0.0,0.0);
        int r = 0;
        int d;
        int ktwiddle = 0;
        int p;
        // Log2 of cols //
        int cols_aux = cols;
        while ( cols_aux >>= 1)
        {
          r++;
        }
        // bitreversal //
        BitReversal_1xN(fft_r,r);
        for(int i=1; i <= r; i++){</pre>
                d = 1 << i;
                ktwiddle = 0;
                for(int k = 0; k < d/2; k++)
                 for(int j = k; j < cols; j+=d){</pre>
                        t = twiddles[ktwiddle]
                                 *fft_r.at<complex<double> >(0,j + d/2);
```

z = fft_r.at<complex<double> >(0,j);
fft_r.at<complex<double> >(0,j) = z + t;
fft_r.at<complex<double> >(0,j + d/2) = z - t;

```
}p = 1 << (r-i);
ktwiddle += p;
}</pre>
```

}

}

1D FFT function for columns.

```
{
          r++;
        }
        // bitreversal //
        BitReversal_Nx1(fft_c,r);
        for(int i=1; i <= r; i++){</pre>
         d = 1 << i;
         ktwiddle = 0;
         for(int k = 0; k < d/2; k++)
          for(int j = k; j < rows; j+=d){</pre>
                   t = twiddles[ktwiddle]
                   *fft_c.at<complex<double> >(j + d/2,0);
                   z = fft_c.at<complex<double> >(j,0);
                   fft_c.at<complex<double> >(j,0) = z + t;
                   fft_c.at<complex<double> >(j + d/2,0) = z - t;
                 }
                 p = 1 << (r-i);
                ktwiddle += p;
         }
 }
}
```

2.3.5 GUI for FFT2D

A graphical User Interface (GUI) has been implemented to improve the execution. In this section we explain how to run this GUI to compute the 2-dimensional FFT for an image.

INPUT IMAGE COMPARE RESULTS
SELECT IMAGE FFT2D FFT2D(matlab)

Figure 2.5: Two dimensional FFT GUI.

We can see in Figure 2.5 the different parts of the GUI:

- Input Image,

If you press *SELECT IMAGE*, then a file dialog will appear where you can choose the input image to be processed.


- Transforms,

We can compute the *Matlab* FFT and our algorithm. The button FFTSHIFT is used to centrate the result, and then you can see the simetry of the FFT.

- TRANSFORMS		
FFT2D	FFT2D(matlab)	FFTSHIFT

- Plot,

There are four different plots. The real part, the imaginary part, the phase, which represent the angle, and the default plot logarithm of the spectrum.

PLOT	
IMAG PART	REAL PART
PHASE	LOG(ABS)

- Time.

Display the time spent in each algoritm in seconds. They appear after the images are plotted.

TIME	
FFT2D(matlab) (s)	
FFT2D (s)	

Example 2.6. In Figure 2.6 we compute the FFT of a RGB image and in Figure 2.7 we compute the FFT of a gray image, the size of the images is 512×512 . In both results we plot the shifted spectrum of them.

We can also see the time spent in Matlab and in our algorithm.



Figure 2.6: Lena.png.

Ast_Fourier_Transform2D		
- INPUT IMAGE-	COMPARE RESULTS	
FTZD FTZD(restilab)	FFT2D FFT2D FFT2D FFT2D FFT2D FFT2D FFT2D FFT2D FFT2D FFT2D FFT2D FFT2D	lab)(3) 1.1859 (3) 2.0074

Figure 2.7: Mandril.png.

Chapter 3

Intensity Transformations and Filtering

Spatial domain techniques operate directly on the pixels of an image. Given an image of dimensions $M \times N$ we have

$$I_2(x,y) = T[I_1(x,y)], (x,y) \in M \times N,$$

where I_1 is the input image, I_2 is the output image and T is the transformation on I_1 .

Definition 3.1. A spatial neighborhood about a point (x, y) is a square or a rectangular region centered at (x, y).

Then, the operator T is applied in each location of the image in the neighborhood selected, to perform the output value of $I_2(x, y)$. See [2].

3.1 Intensity Transformations

The simplest form of the transformation T is when the size of the neighborhood is 1×1 , and is called *intensity transformation* of the image, because the output value depends only on the value of the input.

1. Adjust image function.

This function changes the input range of values of the image by a specified range.

$$T[I(x,y)] = \begin{cases} low_out & \text{if } I(x,y) < low_in \\ low_out + (I(x,y) - low_in)\frac{high_out - low_out}{high_in - low_in} & \text{if } I(x,y) \in C \\ high_out & \text{if } I(x,y) > high_in \end{cases}$$

where $C = [low_in, high_in].$

Here is the implementation of such transformation.

Code _____

```
function out_image = adimage(I,range_in, range_out,gamma)
disp(nargchk(1, 4, nargin));
if nargin == 1
   gamma = 1;
   range_in = [min(min(I)) max(max(I))];
   range_out = [0 1];
elseif nargin == 3
   gamma = 1;
end
[rows, cols, ch] = size(I);
% Check that the ranges are not empty
if isempty(range_in)
   range_in = [0 1];
end
if isempty(range_out)
   range_out = [0 \ 1];
end
low_in = range_in(1);
```

```
high_in = range_in(2);
low_out = range_out(1);
high_out = range_out(2);
%%%%% Transform %%%%%
%% All the c_i are auxiliar matrix %%
c1 = (I \le low_in);
c2 = (I \ge high_in);
c3 = c1|c2;
c4 = (low_in < I < high_in);
out = (I-low_in).*((high_out-low_out)/(high_in-low_in)) + low_out;
c5 = c4.*out;
c6 = ~c3;
out_image = c6.*out + c1.*low_out + c2.*high_out;
out_image = out_image.^gamma;
figure;
subplot(1,2,1);imshow(I);title('Original image');
subplot(1,2,2);imshow(out_image);title('Adjust image');
```

Example 3.2. Execution of the program with differents ranges and images.

- (a) In Figure 3.1 if we get all the possible intensity values as input range and transform it to a smaller one, then the output image has less contrast and the gray values are very similar.
- (b) In Figure 3.2 the output range is reversed, so we obtain its corresponding negative image.
- (c) In Figure 3.3 the contrast of the output image is achieved by selecting the intermediate intensity values and arranging them to [0, 1].



Figure 3.1: Input range: [0 1] & output range: [.3 .7].



Figure 3.2: Input range: $[0\ 1]$ & output range: $[.7\ .3]$.



Figure 3.3: Input range: $[.3\ .7]$ & output range: $[0\ 1].$

2. Logarithmic transformation.

It is defined by

$$T[I(x,y)] = c \cdot \log(1 + double(I(x,y))),$$

where $double(I(x, y)) \in [0, 1]$.

The application of this transform is basically compression of the ranges. For example, it applies when we have to represent the spectrum of the Fourier Transform, whose values are usually high.

Here is the implementation of such transformation.

```
Code
function out = logarithmic(I,c)
disp(nargchk(1, 2, nargin));
if nargin==1
    c = 1;
end
out = c.*(1 + log(I));
figure;
subplot(1,2,1);imshow(I);title('Original image');
subplot(1,2,2);imshow(out);title('Logarithmic image');
```

Example 3.3. Suppose we want to plot the result of the Fast Fourier Transform of an image. The output is a complex image, and we will represent the absolute value of it. The problem is that there are high values that are not represented because of the range. Then, the logarithmic transform achieves to represent in the plot those high values.



Figure 3.4: Logarithmic transformation.

3. Contrast stretching transformation.

It compresses the inputs levels lower than m into a narrow range of dark levels in the output image; and it compresses the values above m into narrow band of light levels in the output. The function for differents inputs of m and E can be seen in Figure 3.6.



Figure 3.5: Constract stretching function

It is defined by means of the transformation:

$$T[I(x,y)] = \frac{1}{1 + (m/I(x,y))^E},$$

where E controls the slope of the function. Here is the implementation of such transformation.

```
Code
function out=contrast_stretching(I,m,E)
disp(nargchk(1, 3, nargin));
if nargin == 1
    m = (max(max(I))-min(min(I)))/2;
    E = 20;
end
out = 1./(1 + (m./I).^E);
figure;
subplot(1,2,1);imshow(I);title('Original image');
subplot(1,2,2);imshow(out);title('Contrast&Stretching image');
```

Example 3.4. In the example below, we apply the transform to the image estatua.jpg. The values of $m \in [0, 1]$ because this image takes values in this range.



Figure 3.6: Constract stretching transformation

3.1.1 Histogram and Equalization

Histograms are used in image processing for enhancement, compression, segmentation, etc. The histogram h computes how many times an intensity value is repeated:

$$h(r_k) = n_k$$

where r_k is the k intensity value and n_k is its corresponding absolute frequency.

When we calculate the histogram of an image it is necessary to specify how many bins (adjacent rectangles over discrete intervals) will be used; by default, the number of bins will be 256.

Equalization of the histogram is a transformation that generates an image whose intensity levels are basically equal and then all intensity value range is covered. There are many ways to compute the equalization and we will mostly consider the following:

Define

$$p_r(r_k) = \frac{n_k}{N},$$

with N all the intensity values, as the probability density function (PDF). Then the transformation is given by

$$s_k = \sum_{j=1}^k p_r(r_j),$$

which is called the *cumulative distribution function (CDF)*.

Remark 3.5. In the process of histogram equalization, information data is lost. So this transformation is useful for viewing the image better, but it is not convenient to use the output image for image processing later.

The histogram function has been implemented for RGB, red-green-blue channels, color images, and gray-level images. Here is the code for each function and histogram equalization. Histogram function of a gray-level image.

```
Code -
function [h,x]=histdouble(I,bin)
disp(nargchk(1, 2, nargin));
if nargin == 1
    bin = 256;
end
if isa(I,'double')
    A=1;
    x = 0:1/255:1;
elseif isa(I,'uint8')
    A=255;
    x = 0:255;
end
[rows cols] = size(I);
h = zeros(1, 256);
for i=1:rows
    for j=1:cols
        p = floor(I(i,j).*((bin-1)/A) + 1.5);
        if p>256
            p=256;
        end
        h(p) = h(p) + 1;
    end
\operatorname{end}
```

Histogram function of a RGB image.

Image equalization.

_____ Code _____

```
function [im_eq,h_eq,s,h1]=histogram_eq(I,h,n)

disp(nargchk(1, 3, nargin));
if nargin == 2
    n = 256;
end
if isa(I,'double')
    A=1;
elseif isa(I,'uint8')
    A=255;
end
[rows cols] = size(I);

s = zeros(1,n);
for k=1:n
    s(k) = sum(h(1:k));
```

end

```
im_eq = (s(floor(I.*((n-1)/A)+1.5))-min(s))./(rows*cols-min(s));
h_eq = histdouble(im_eq);
h1 = zeros(1,n);
for k=1:n
    h1(k) = sum(h_eq(1:k));
end
```

GUI of histogram equalization

In this section we explain how to run this GUI, enumerating its parts:

- Input Image.
- Histogram and equalization.
- Plot histograms.
- PDF and CDF plots.
- Equalize image.

If you press *SELECT IMAGE*, then a file dialog will appear where you can choose the input image to be processed.



Each button compute the histogram and the equalized histogram and plot them.

Histogram	
Equalize hist.	

There are 3 different ways to plot the histograms: with bars, stems or with a simple plot. The menu list of RGB-channel is only for images of 3 channels, then you can choose which one will be plot.

Plots
Bar Stem
Plot
RGB-channel

Check the corresponding box to plot the PDF and CDF defined as above.

Curves	Ī
Show PDF	
Show CDF	

Equalize image

This button displays the equalized image.

Example 3.6. Let us consider lena.png and estatua.jpg to show the image equalization of images of different number of channels. In Figure 3.7 we compute it for a gray-level image, and in Figure 3.8 we compute it for a color image.



Figure 3.7: Histogram of a gray-level image.



Figure 3.8: Histogram of an RGB image.

3.2 Filtering

3.2.1 Spatial Filter Domain

Definition 3.7. Given a neighborhood as defined in last section, we define linear spatial filtering as the multiplication of each pixel in the neighborhood by a corresponding coefficient and summing the results to obtain a new value at pixel (x, y).

The coefficients are arranged as a matrix called *mask*, *filter*, *filter mask*, *kernel*, *template* or *window*. The most common are the first three. Figure 3.9 shows how to process the mask filter accross the image.



Figure 3.9: Filtering image

Definition 3.8. 1. Correlation is a process of passing the mask w by the image.

2. Convolution is a process of passing the mask w rotated 180° by the image.

Example 3.9. Consider a one dimensional signal $f = 0\ 0\ 0\ 1\ 0\ 0\ 0\ 0$ and a mask $w = 1\ 2\ 3\ 2\ 0$.

1. Correlation

```
First shift

1 \ 2 \ 3 \ 2 \ 0

Second shift

1 \ 2 \ 3 \ 2 \ 0

Fourth shift

1 \ 2 \ 3 \ 2 \ 0

Fourth shift

1 \ 2 \ 3 \ 2 \ 0

Fourth shift

1 \ 2 \ 3 \ 2 \ 0

Final position

0 \ 0 \ 0 \ 1 \ 0 \ 0 \ 0 \ 0

Result 0 \ 0 \ 0 \ 2 \ 3 \ 2 \ 1 \ 0 \ 0 \ 0
```

2. Convolution

Result 0 0 0 1 2 3 2 0 0 0 0 0

3.2.2 Boundary options in the filter

We have to extend the image boundaries because when we apply the filter, we put the filter center c in each pixel image. It can be seen in Figure 3.10.

PADDING: The boundaries are extended by the value p. In our case, the padding is made with zeros.

REPLICATE: The size of the image is extended by replicating its values in the borders.

REFLEXION: The size of the image is extended by symmetry followed by a mirrorreflecting across its borders.



Figure 3.10: Image boundaries

3.2.3 Spatial filters

AVERAGE: A rectangular averaging filter of dimension $r \times c$, where r is the number of rows and c the number of columns.

DISK: A circular averaging filter of dimension $(2r + 1) \times (2r + 1)$, where r is the radius.

GAUSSIAN: A Gaussian filter of size $r \times c$ and standard deviation $\sigma \geq 0$.

LAPLACIAN: Laplacian filter of size 3×3 and $\alpha \in [0, 1]$.

LOG: Laplacian of Gaussian (LoG) filter of size $r \times c$ and standard deviation $\sigma \geq 0$.

PREWITT: Prewitt mask of size 3×3 that approximates a vertical gradient.

SOBEL: Sobel mask of size 3×3 that approximates an horitzontal gradient.

In Figure 3.11 and 3.12 the Gaussian, Laplacian, LoG and Prewitt, Sobel spatial filters are represented respectively.



Figure 3.11: Gaussian, Laplacian and LoG Filters.



Figure 3.12: Gradient Filters.

Here is the code of the convolution in the spatial domain of an image.

___ Code ____

function conv2D = conv2D(image,filter,method)

```
% Convert image to double (if the image is double, it doesn't change)
image = im2double(image);
```

```
[rows,cols]=size(image);
[rows_f,cols_f]=size(filter);
r = floor(rows_f/2);
c = floor(cols_f/2);
newimage = zeros(rows+2*r,cols+2*c);
newimage(r+1:end-r,c+1:end-c) = image;
if (strcmp(method,'period'))
    % Fill newimage
    newimage(1:r,:) =
    [image(end-r+1:end,end-c+1:end)
    image(end-r+1:end,:)
    image(end-r+1:end,1:c)];
    newimage(end-r+1:end,:) =
    [image(1:r,end-c+1:end) image(1:r,:) image(1:r,1:c)];
    newimage(r+1:end-r,1:c) = image(:,end-c+1:end);
    newimage(r+1:end-r,end-c+1:end) = image(:,1:c);
    % Convolution
    for i=1:rows
        for j=1:cols
           conv2D(i,j) = sum(sum(newimage(i:i+2*r,j:j+2*c).*filter));
        end
    end
elseif (strcmp(method, 'reflexion'))
    % Fill newimage
    newimage(1:r,:) =
    [image(r:-1:1,end:-1:end-c+1)
    image(r:-1:1,:)
    image(r:-1:1,end:-1:end-c+1)];
```

```
newimage(end-r+1:end,:) =
    [image(end:-1:end-r+1,c:-1:1)
    image(end:-1:end-r+1,:)
    image(end:-1:end-r+1,end:-1:end-c+1)];
    newimage(r+1:end-r,1:c) = image(:,c:-1:1);
    newimage(r+1:end-r,end-c+1:end) = image(:,end:-1:end-c+1);
    % Convolution
    for i=1:rows
        for j=1:cols
            conv2D(i,j) = sum(sum(newimage(i:i+2*r,j:j+2*c).*filter));
        end
    end
elseif(strcmp(method,'padding'))
    % Convolution
    for i=1:rows
        for j=1:cols
            conv2D(i,j) = sum(sum(newimage(i:i+2*r,j:j+2*c).*filter));
        end
    end
else
    'method = period,reflexion,padding'
    return
end
% figure;
% subplot(1,2,1);imshow(image,[]);title('Image');
% subplot(1,2,2);imshow(conv2D,[]);title('Convolution 2D');
```

3.2.4 Frequency Domain

In the frequency domain we use Theorem 1.9.

$$f(x,y) * h(x,y) \iff \hat{f}(u,v) \cdot \hat{h}(u,v),$$

where f is the image and h is the filter.

- **Remark 3.10.** 1. If filter size is not equal to image size, then padd with zeros to achieve the same size.
 - 2. First we compute the convolution in the frequency domain and then we apply the IFFT to the result.

Here is the code for the convolution in the frequency domain.

```
_ Code _
function conv2D = conv2D_f(I,filter)
%%
%%
  Convolution in the frequency domain
%%
[rows cols] = size(I);
[rows_f cols_f] = size(filter);
rows_pad = padding1D(rows);
cols_pad = padding1D(cols);
I_pad = zeros(rows_pad,cols_pad);
I_pad(1:rows,1:cols) = I;
filter_pad = zeros(rows_pad,cols_pad);
filter_pad(1:rows_f,1:cols_f) = filter;
```

```
[fft_image_r fft_image_i] = FFT2_mex.FFT(I_pad);
[fft_filter_r fft_filter_i] = FFT2_mex.FFT(filter_pad);
fft_image = complex(fft_image_r,fft_image_i);
fft_filter = complex(fft_filter_r,fft_filter_i);
conv2D_f = fft_image.*fft_filter;
[conv2D_r conv2D_i] = FFT2_mex.IFFT(conv2D_f);
conv2D1 = complex(conv2D_r,conv2D_i);
conv2D = conv2D1(floor(rows_f/2):rows+floor(rows_f/2),
floor(cols_f/2):cols+floor(cols_f/2));
```

3.2.5 GUI Image Filtering

In this section we explain how to run this *GUI* enumerating its parts:

- Input Image,
- Choose parameters of the filter,
- Plot the filtered image.

If you press *SELECT IMAGE*, then a file dialog will appear where you can choose the input image to be processed.

INPLIT IMAGE		
NU OT NU IOL		
SELECT MAGE		

Introduce all the parameters for the filter and select it (in this order). Then press *Apply* button to compute the transformations.

Average rows cols
Disk rows
Gaussian rows cols sig
Laplacian alpha
LoG rows cols sig
Prewitt 3x3 prewitt filter
Sobel 3x3 sobel filter
Apply

	Select Filter Do
Choose spatial or frequency	
domain.	Frequency Do

Select Filter Domain		
Trequency Domain	Spatial Filtering	

Example 3.11. We have added 'salt and pepper'noise (randomly white and black pixels) in the picture of lena.png. In Figure 3.13 and Figure 3.14 we can observe that the image has been improved, because these filters are called low-pass filters. But in Figure 3.15 the noise is not erased, because it is called a high-pass filter.



Figure 3.13: Average filter.



Figure 3.14: Gaussian filter.



Figure 3.15: Laplacian of Gaussian filter.

Chapter 4

Morphological Image Processing

In this chapter we will introduce morphological operations in image processing, like dilation, erosion, openning, closing... See [2]. In the two first sections we will deal with binary images, and in the last section all the functions will be applied in gray level images.

4.1 Dilation and Erosion

First, we will introduce some correspondences between set theory and its equivalent with binary images.

Definition 4.1. Reflexion of a set A is defined as:

$$\hat{A} = \{ w \mid w = -a, \ a \in A \}.$$

Traslation of a set A is defined by:

$$(A)_{z} = \{c \mid c = a + z, a \in A\}.$$

In table 4.1 we represent the equivalence between set theory and *Matlab* syntax.

Set theory	Matlab
$A \cap B$	A&B
$A \cup B$	A B
A^c	-A
A - B	A& - B

Table 4.1: Set theory & Matlab syntax

Example 4.2. Given two binary images, A and B, calculate the intersection, union, complement and set difference.



4.1.1 Dilation

Definition 4.3. An structuring element is computationally represented as a matrix of 0s and 1s and by an origin.

Definition 4.4. Dilation is an operation that 'grows'or 'thickens'objects in binary images.

$$A \oplus B = \left\{ z \mid (\hat{B})_z \cap A \neq \emptyset \right\}.$$

Dilation is a commutative operation $A \oplus B = B \oplus A$.

Example 4.5. Given a binary image and an structuring element

$$\begin{pmatrix} 1 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 1 & 1 \end{pmatrix}$$

we calculate its dilation:



Figure 4.2: Dilation operation.

We can observe that the output image has increased the size of the rice. That is because dilation selects all those pixels that have intersection with the given structuring element.

We have implemented a function *structel* to calculate the most used structuring elements.

Syntax	Description
structel('diamond', R)	Diamond form of radius R
structel('disk', R)	Disk of radius R
structel('octagon', R)	Octagon of radius R
structel ('rectangle',[$M\ N$])	Rectangle of size $M \times N$
structel('line',L)	Line of length L
structel('square', M)	Square of size $M \times M$

4.1.2 Erosion

Definition 4.6. Erosion 'shrinks'or 'thins' objects in a binary image. It is defined as:

$$A \ominus B = \{ z \mid (B)_z \cap A^c \neq \emptyset \}$$

Example 4.7. With the same binary image and structuring element (4.5) used in dilation we will calculate its erosion:



Figure 4.3: Erosion operation.

The output image has decreased the rice size, because erosion deletes all those pixels that have not the same neighborhood as the given structuring element.

4.1.3 Implementation in Matlab

Dilation

```
_ Code ___
 function A2 = dilate(A,B)
[rows,cols]=size(A);
[rowsB,colsB] = size(B);
% Calcule the center coordinates of B
cx = ceil(colsB/2);
cy = ceil(rowsB/2);
% Calculate new image
new_image = zeros(rows+rowsB-1,cols+colsB-1);
if mod(colsB,2)~=0
    lastx = cols+colsB-1 -cx+1;
    deltax=cx-1;
else
    lastx = cols+colsB-1 -cx;
    deltax=cx;
end
if mod(rowsB,2)~=0
    lasty = rows+rowsB-1 -cy+1;
    deltay=cy-1;
else
    lasty = rows+rowsB-1 -cy;
    deltay= cy;
end
new_image(cy:lasty,cx:lastx) = A;
%% Calcule reflection of B
R_B = B(end:-1:1, end:-1:1);
%Calculate dilation
```

```
A2 = zeros(rows,cols);
for i=cy:lasty
    for j=cx:lastx
        a = sum(sum(new_image(i-cy+1:i+deltay,j-cx+1:j+deltax).*R_B));
        if a~=0
            A2(i-cy+1,j-cx+1)=1;
        end
    end
end
    Erosion
                               Code _____
  function A2=erode(A,B)
[rows,cols]=size(A);
[rowsB,colsB] = size(B);
\% Calcule the center coordinates of B
cx = ceil(colsB/2);
cy = ceil(rowsB/2);
% Calculate new image
new_image = zeros(rows+rowsB-1,cols+colsB-1);
if mod(colsB,2)~=0
    lastx = cols+colsB-1 -cx+1;
else
    lastx = cols+colsB-1 -cx;
end
if mod(rowsB,2)~=0
    lasty = rows+rowsB-1 -cy+1;
else
    lasty = rows+rowsB-1 -cy;
end
```

```
new_image(cy:lasty,cx:lastx) = A;
%Calculate
sumel=sum(sum(B));
A2 = zeros(rows,cols);
for i=cy:lasty
    for j=cx:lastx
        a = sum(sum(new_image(i-cy+1:i+cy-1,j-cx+1:j+cx-1).*B));
        if a == sumel
            A2(i-cy+1,j-cx+1) = 1;
        end
    end
end
```

4.2 Opening, Closing and Hit-or-Miss

In practice, in image processing we combine dilations and erosions with different structuring elements. We enhance the most common combinations: opening, closing and hit-or-miss.

4.2.1 Opening

Definition 4.8. The morphological operation openning of the image A by the element B is denoted by $A \circ B$. It is defined as:

$$A \circ B = (A \ominus B) \oplus B.$$

This operation eliminates the regions that do not contain the structuring element, smooth the countours and break weak connections between objects.

4.2.2 Closing

Definition 4.9. The morphological operation of closing is defined as:

$$A \bullet B = (A \oplus B) \ominus B$$

This operation smooths the countours, and fill the holes that are smaller than the structuring elements.

Example 4.10. Given a fingerprint image with noise, first we apply opening to delete noisy dots. But this operation produces holes in the fingerprint. Then, if we apply the closing operation in this output image we obtain the desire result.



(a) Fingerprint with noise



(c) Opening followed by Closing

Figure 4.4: Opening and Closing.

(b) Opening

4.2.3 Hit-or-Miss

This transformation is useful for identifying configurations of pixels like isolated pixels and endings of lines.

Definition 4.11. The Hit-or-Miss of the set A by the element $B = (B_1, B_2)$ is defined as:

$$A \otimes B = (A \ominus B_1) \cap (A^c \ominus B_2).$$

The *Hit-or-Miss* consists of all pixels which are in B_1 (*hit*), and those that have not pixels in B_2 (*miss*).

4.2.4 Implementation in Matlab

We only present the implementation of *Hit-or-Miss* because *opening* and *closing* operations are form using *dilations* and *erosions*.

Hit-or-Miss

```
Code _____
function [A2,C2]=HitorMiss(A,B1,B2)
%%
%%
      Hit-or-Miss transformation
%%
      Description:
%%
      Structuring elements (B1,B2)
%%
      HoM:=erode(A,B1) intersection erode(Ac B2)
%%
%% In general, supose that B has odd dimension
A=im2double(A);
%% Image A must be a binary image %
A = im2bw(A);
% Calculate A supplementary
A_c = ~A;
C1 = erode(A, B1);
C2 = erode(A_c, B2);
% Calculate the intersection between C1 & C2:
A2 = C1\&C2;
```

We introduce *bottom-hat* and *top-hat*. These are operations that extract small elements and details from a given image.

Bottom-hat returns the difference between its closing by a structuring element and the image.

```
A2 = imsubtract(A1,A);
```

Top-hat returns the difference between the image and its opening by some structuring element.

```
A2 = imsubtract(A1,A);
```
4.2.5 Lookup Tables

A lookup table is a data structure used to replace a runtime computation with a simpler array indexing operation. In image processing the nomenclature is LUT.

The technique is to precompute the output pixel value for every possible neighborhood configuration and then store the answer in a table for later use. For example, we have $2^9 = 512$ different 3×3 pixels configurations in a binary image.

To make the LUT we have to assign a unique value to each configuration, we will multiply each element of our neighborhood with the following matrix

1	8	64
2	16	128
4	32	256

Example 4.12. Let us consider the 3×3 neighborhood

1	1	0
1	0	1
1	0	1

the value that will represent this matrix uniquely is perform as follows: 1(1) + 2(1) + 4(1) + 8(1) + 16(0) + 32(0) + 64(0) + 128(1) + 256(1) = 399.

We will implement two functions: *makelookuptab* and *applylookuptab*. The first one constructs a lookup table and the last one processes the binary images using this lookup table.

Example 4.13. The lookup table is contructed next by calling makelut with a function handle to conwaylaws



Figure 4.5

4.2.6 Implementation in Matlab

Auxiliar function *matrix3x3*. Here, we create all the 3×3 matrices of 1's and 0's, there are 2^9 different matrices.

Auxialiar function sum1. It calculates the value for each matrix in the above function matrix3x3.

Make lookup table.

```
_ Code _____
```

```
return
end
if n==2
    lut = zeros(pow2(4),2,2,2);
   m = matrix2x2();
    for i=1:pow2(4)
        lut(i,1) = handle(reshape(m(i,:),[2,2]));
        lut(i,2,:,:) = reshape(m(i,:),[2,2]);
    end
elseif n==3
    lut = zeros(pow2(9),1);
    m = matrix3x3( );
   for i=1:pow2(9)
        lut(i) = handle(reshape(m(i,:),[3,3]));
    end
else
    'Second argument must be 2 or 3'
    return
end
```

Apply lookup table.

_____ Code _____

```
function fnew=applylookuptab(f,lut)
```

```
[r,c]=size(lut);
aux = f;
if r == 16
    m = matrix2x2();
    size(m)
```

```
for i=1:r
    if lut(i)==1
        fnew = mask(aux,reshape(m(i,:),[2,2]));
        aux = fnew;
    end
end
else
    m = matrix3x3();
    for i=1:r
        if lut(i)
            fnew = mask(aux,reshape(m(i,:),[3,3]));
            aux = fnew;
        end
end
end
end
```

4.3 Morphological operations in gray images

All the binary morphological operations, except Hit-or-Miss transform, have extensions to gray images. In this section, we will work with dilation, erosion, opening and closing. Finally we will introduce how to obtain the gradient of the image with some of these operations.

4.3.1 Dilation and Erosion

Definition 4.14. The gray-scale dilation of an image by a structuring element, is denoted by

$$(I \oplus B)(x, y) = max \left\{ I\left(x - x', y - y'\right) + B\left(x', y'\right) | (x', y') \in B \right\},\$$

where I is the image and B is the structuring element. We assume that the value $I(x,y) = -\infty$ outside the image.



Figure 4.6: Dilation in gray images

```
Code _____
function A2 = rkl_dilate_gray(A,B)
[rows,cols]=size(A);
[rowsB,colsB] = size(B);
% Calcule the center coordinates of B
cx = ceil(colsB/2);
cy = ceil(rowsB/2);
% Calculate new image
new_image = zeros(rows+rowsB-1,cols+colsB-1);
new_image(:) = -Inf;
if mod(colsB,2)~=0
    lastx = cols+colsB-1 -cx+1;
    deltax=cx-1;
else
    lastx = cols+colsB-1 -cx;
    deltax=cx;
end
```

```
if mod(rowsB,2)~=0
    lasty = rows+rowsB-1 -cy+1;
    deltay=cy-1;
else
    lasty = rows+rowsB-1 -cy;
    deltay= cy;
end
new_image(cy:lasty,cx:lastx) = A;
%% Calcule reflection of B
R_B = B(end:-1:1, end:-1:1);
%Calculate dilation
A2 = zeros(rows,cols);
for i=cy:lasty
    for j=cx:lastx
        a = new_image(i-cy+1:i+deltay,j-cx+1:j+deltax).*R_B;
        A2(i-cy+1, j-cx+1)=max(max(a));
    end
end
```

Definition 4.15. The gray-scale erosion of an image by a structuring element, is denoted by

$$(I \ominus B)(x, y) = min\left\{I\left(x + x^{'}, y + y^{'}\right) - B\left(x^{'}, y^{'}\right) | (x^{'}, y^{'}) \in B\right\},\$$

where I is the image and B is the structuring element. We assume that the value $I(x, y) = +\infty$ outside the image.



Figure 4.7: Erosion in gray images.

```
Code _____
function A2 = rkl_erosion_gray(A,B)
[rows,cols]=size(A);
[rowsB,colsB] = size(B);
% Calcule the center coordinates of B
cx = ceil(colsB/2);
cy = ceil(rowsB/2);
% Calculate new image
new_image = zeros(rows+rowsB-1,cols+colsB-1);
new_image(:,:) = +Inf;
if mod(colsB,2)~=0
    lastx = cols+colsB-1 -cx+1;
else
    lastx = cols+colsB-1 -cx;
end
if mod(rowsB,2)~=0
    lasty = rows+rowsB-1 -cy+1;
else
```

```
lasty = rows+rowsB-1 -cy;
end
new_image(cy:lasty,cx:lastx) = A;
A2 = zeros(rows,cols);
for i=cy:lasty
    for j=cx:lastx
        a = new_image(i-cy+1:i+cy-1,j-cx+1:j+cx-1).*B;
        a(a==0) = +Inf;
        A2(i-cy+1,j-cx+1) = min(min(a));
    end
end
```

4.3.2 Opening and Closing

Definition 4.16. The opening of an image by a structuring element is denoted as

$$I \circ B = (I \ominus B) \oplus B,$$

where I is the image, B the structuring element and the operations \oplus, \ominus as defined above.



Figure 4.8: Opening in gray images.

Definition 4.17. The closing of an image by a structuring element is denoted as

$$I \bullet B = (I \oplus B) \ominus B,$$

where I is the image, B the structuring element and the operations \oplus, \ominus as defined above.



Figure 4.9: Closing in gray images.

4.3.3 Applications: bottom-hat, top-hat and morphological gradient

Definition 4.18. The Bottom-hat transformation is defined as

 $I \bullet B - I$,

where B is the structuring element, and I is the image.



Figure 4.10: Bottom-hat of estatua.jpg

A2 = imsubtract(A1,A);

Definition 4.19. The Top-hat transformation is defined as

 $I - I \circ B$,

where B is the structuring element, and I is the image.



Figure 4.11: Top-hat of *estatua.jpg*

Definition 4.20. The morphological gradient transformation is defined as

 $I \oplus B - I \ominus B,$

where B is the structuring element, and I is the image.



Figure 4.12: Morphological gradient of estatua.jpg.

4.3.4 GUI Morphological Operations

In this section we explain how to run the two GUI's of morphological operations, one for binary images and the other for intensity images.

Morphological operations GUI for binary images.

BW_morphological		
MORPHOLOGICAL OPERATIONS		
Lookup table Pundon Egi conweytewe Apply		
Citlers	Copen marger	
	Satinga	

Figure 4.13: Binary images.

- Open the image to be processed.

- Choose the morphological operation. It will appear another window to choose the structuring element. Then press *Apply* button. See Figure 4.15.
- The output image will appear in the main window.

Morphological operations GUI for intensity images.



Figure 4.14: Intensity images.

- Open the image to be processed.
- Choose the morphological operation. It will appear another window to choose the structuring element. Then press *Apply* button. See Figure 4.15.
- The output image will appear in the main window.

There are two windows to choose the structuring element. The one that has two panels is only used for *Hit-or-Miss* transformation.

Structuring element		
MATRIX		
	E.g: [1 1 1;0 1 0;1 1 1]	
Predefined		
	OCTAGON	
R: radius	R: radius	
) DISK	RECTANGLE	
R: radius	[rows; columns]	
	SQUARE	
[length; degrees]	N: size	

inst Structuring Element		Second Structuring Benerit	
MATRIX		MATRIX	
	E.gr. [111,010;111]		Eg[111;010;111]
Predefined		Predefined	
C DIAMOND	O OCTAGON	© DAMOND	OCTAGON
R radius	R radius	R redus	R: radius
O DEX		© DEK	C RECTANGLE
R raduz	(rows; columns)	R: radus	(nowa; columna)
O LNE	C SQUARE	© LNE	C SQUARE
[ength; degrees]	N. size	(length; degrees)	N size

(a) One structuring element.

(b) Two structuring elements.

Figure 4.15: Structuring element windows.

Chapter 5

Wavelets on \mathbb{Z}_N

5.1 Introduction

In this section we introduce the motivation of *Wavelets* analysis in front of *Fourier* analysis. We will continue working with discrete signals in $\ell^2(\mathbb{Z}_N)$. See [4].

Definition 5.1. We say that a vector $z \in \ell^2(\mathbb{Z}_N)$ is localized in space near $n_0 \in \mathbb{Z}_N$, if most of the components z(n) of z are 0 or al least relatively small, except for a few values of n close to n_0 .

A Fourier basis element is not localized, because it has the same magnitude $\frac{1}{\sqrt{N}}$ for every $n \in \mathbb{Z}_N$. See 1.1.

Now, suppose $B = \{v_0, \ldots, v_{N-1}\}$ is a basis for $\ell^2(\mathbb{Z}_N)$ such that all the basis elements of B are *localized in space*. For a vector z, we can write $z = \sum_{n=0}^{N-1} a_n v_n$, $a_n \in \mathbb{R}$. Suppose that we wish to focus on the position of z near some particular point n_0 . Terms involving basis vectors that are 0 or relatively small near n_0 can be deleted in without changing the behaviour. Then we can replace a sum over Nterms by a smaller sum.

5.2 Wavelets on \mathbb{Z}_N

Notation 5.2. We introduce new notation, and we recall the notation used above.

1.
$$\tilde{w}(n) = \overline{w(-n)} = \overline{w(N-n)}.$$

$$2. \ (\tilde{w})(n) = \hat{w}(n).$$

3.
$$(R_k z)(n) = z(n-k)$$
.

4.
$$(z * \tilde{w})(k) = \langle z, R_k w \rangle.$$

We call \tilde{w} the conjugate reflection of w.

Lemma 5.3. Suppose $z, w \in \ell^2(\mathbb{Z}_N)$. For any $k \in \mathbb{Z}$

$$z * \tilde{w}(k) = \langle z, R_k w \rangle,$$
$$z * w(k) = \langle z, R_k \tilde{w} \rangle.$$

Remark 5.4. $\widetilde{(\tilde{w})} = w$: $\widetilde{(\tilde{w})} = (\overline{w(-n)}) = \overline{(\overline{w(n)})} = w(n).$

Proof. By definition and commutativity,

1.

$$\langle z, R_k w \rangle = \sum_{n=0}^{N-1} z(n) \overline{R_k w(n)} = \sum_{n=0}^{N-1} z(n) \overline{w(n-k)}$$
$$= \sum_{n=0}^{N-1} z(n) \tilde{w}(k-n) = \tilde{w} * z(k) = z * \tilde{w}(k).$$

2.

$$\langle z, R_k \tilde{w} \rangle = \sum_{n=0}^{N-1} z(n) \overline{R_k \tilde{w}(n)} = \sum_{n=0}^{N-1} z(n) \overline{\tilde{w}(n-k)}$$
$$= \sum_{n=0}^{N-1} z(n) \widetilde{\tilde{w}}(k-n) = w * z(k) = z * w(k).$$

Lemma 5.5. Let $w \in \ell^2(\mathbb{Z}_N)$. Then $\{R_k w\}_{k=0}^{N-1}$ is an orthonormal basis for $\ell^2(\mathbb{Z}_N)$ if and only if $|\hat{w}(n)| = 1$ for all $n \in \mathbb{Z}_N$.

Proof. Let us introduce the Dirac Delta distribution,

$$\delta(n) = \begin{cases} 1 & \text{if } n = 0 \\ 0 & \text{if } n = 1, \dots, N - 1. \end{cases}$$

By a simple calculation, we obtain $\hat{\delta}(n) = 1$ for all $n \in N$:

$$\hat{\delta}(n) = \sum_{m=0}^{N-1} \delta(m) e^{-2\pi i m n/N} = 1.$$

 $\{R_k w\}_{k=0}^{N-1}$ is an orthonormal basis if and only if

$$\langle w, R_k w \rangle = \begin{cases} 0 & \forall k = 1, \dots, N-1 \\ 1 & k = 0 \end{cases}$$

By notation,

$$(w * \tilde{w})(k) = \delta(k) \iff (w * \tilde{w})(n) = \hat{\delta}(n) = 1 \ \forall n$$
$$\iff \hat{w}(n)(\tilde{w})(n) = 1 \ \forall n \iff |\hat{w}(n)| = 1.$$

Definition 5.6. Suppose N is an even integer, N = 2M for some $M \in \mathbb{N}$. An orthonormal basis for $\ell^2(\mathbb{Z}_N)$ of the form

$$\{R_{2k}u\}_{k=0}^{M-1} \bigcup \{R_{2k}v\}_{k=0}^{M-1}$$

for some $u, v \in \ell^2(\mathbb{Z}_N)$ is called a first-stage wavelet basis.

Our goal is to determine when a pair u, v generates a *first-stage wavelet*. First we will introduce the necessary background for the theorem that will determine that pair.

•

Lemma 5.7. Suppose $M \in \mathbb{N}$, N = 2M, $z \in \ell^2(\mathbb{Z}_N)$. Define $z^* \in \ell^2(\mathbb{Z}_N)$ by

$$z^*(n) = (-1)^n z(n), \text{ for all } n.$$

Then,

$$\widehat{(z^*)}(n) = \hat{z}(n+M), \text{ for all } n.$$

Proof. By definition,

$$\widehat{(z^*)}(n) = \sum_{k=0}^{N-1} z^*(k) e^{-2\pi i k n/N} = \sum_{k=0}^{N-1} (-1)^k z(k) e^{-2\pi i k n/N}$$
$$= \sum_{k=0}^{N-1} z(k) e^{-i\pi k} e^{-2\pi i k n/N} \underbrace{=}_{\frac{-2\pi i k n}{N} - \pi i k = \frac{-2\pi i k (n+M)}{N}} \sum_{k=0}^{N-1} z(k) e^{-2\pi i k (n+M)/N}$$
$$= \widehat{z}(n+M).$$

Remark 5.8. For any $z \in \ell^2(\mathbb{Z}_N)$, with N even

$$(z+z^*)(n) = z(n)(1+(-1)^n) = \begin{cases} 2z(n) & \text{if } n \text{ is even,} \\ 0 & \text{if } n \text{ is odd.} \end{cases}$$

Lemma 5.9. Suppose $M \in \mathbb{N}$, N = 2M and $w \in \ell^2(\mathbb{Z}_N)$. Then $\{R_{2k}w\}_{k=0}^{M-1}$ is an orthonormal set with M elements if and only if

$$|\hat{w}(n)|^2 + |\hat{w}(n+M)|^2 = 2$$
 for $n = 0, \dots, M-1$

Proof. $\{R_{2k}w\}_{k=0}^{M-1}$ is an orthonormal basis if and only if

$$\langle w, R_{2k}w \rangle = \begin{cases} 1 & k = 0\\ 0 & k = 1, \dots, M - 1 \end{cases}$$
$$(w * \tilde{w})(n) + (w * \tilde{w})^*(n) = \begin{cases} 2 & n = 0\\ 0 & n = 1, \dots, N - 1 \end{cases}$$

$$(w * \tilde{w}) + (w * \tilde{w})^* = 2\delta \iff$$
$$(w * \tilde{w})(n) + [(w * \tilde{w})^*](n) = 2\hat{\delta} = 2 \forall n \iff$$
$$\underset{Lemma \ 5.7, \ Notation \ 5.2}{\longleftrightarrow}$$
$$|\hat{w}(n)|^2 + |\hat{w}(n+M)|^2 = 2 \forall n.$$

Definition 5.10. Suppose $M \in \mathbb{N}$, N = 2M and $u, v \in \ell^2(\mathbb{Z}_N)$. For $n \in \mathbb{Z}$, define A(n) the system matrix of u and v by

$$A(n) = \frac{1}{\sqrt{2}} \begin{pmatrix} \hat{u}(n) & \hat{v}(n) \\ \hat{u}(n+M) & \hat{v}(n+M) \end{pmatrix}$$

Now, we can introduce the necessary and sufficient conditions to determine *first-stage wavelet basis* pair.

Theorem 5.11. Suppose $M \in \mathbb{N}$, N = 2M. Let $u, v \in \ell^2(\mathbb{Z}_N)$. Then,

$$B = \{R_{2k}u\}_{k=0}^{M-1} \bigcup \{R_{2k}u\}_{k=0}^{M-1}$$

= $\{u, R_2u, R_4u, \dots, R_{N-2}u, v, R_2v, R_4v, \dots, R_{N-2}v\}$

is an orthonormal basis for $\ell^2(\mathbb{Z}_N)$ if and only if A(n) is unitary for each $n = 0, \ldots, M - 1$.

Equivalently, B is a first-stage wavelet basis for $\ell^2(\mathbb{Z}_N)$ if and only if

1. $|\hat{u}(n)|^2 + |\hat{u}(n+M)|^2 = 2 \ \forall n,$ 2. $|\hat{v}(n)|^2 + |\hat{v}(n+M)|^2 = 2 \ \forall n,$ 3. $\hat{u}(n)\overline{\hat{v}(n)} + \hat{u}(n+M)\overline{\hat{v}(n+M)} = 0 \ \forall n.$

Proof. We know that $\{R_{2k}u\}_{k=0}^{M-1}$ is an orthonormal set if and only if first condition is satisfied.

The same for $\{R_{2k}v\}_{k=0}^{M-1}$ with the second condition.

So, we only have to proof the last one:

$$\langle R_{2k}u, R_{2k}v \rangle = 0, \ \forall j \neq k \iff (3) \ is \ satisfied.$$

$$(u * \tilde{v})(2k) = \langle u, R_{2k}v \rangle = 0 \iff$$
$$(u * \tilde{v})(2k) = 0, \ \forall k = 0, \dots, M - 1 \iff$$
$$(u * \tilde{v})(n) + (u * \tilde{v})^*(n) = 0 \iff$$
$$(u * \tilde{v})(n) + [(u * \tilde{v})^*(n)] = 0 \iff$$
$$\hat{u}(n)\overline{\hat{v}(n)} + \hat{u}(n + M)\overline{\hat{v}(n + M)} = 0 \ \forall n.$$

Example 5.12. 1. Haar first-stage wavelet basis.

$$u = \left(\frac{1}{\sqrt{2}}, \frac{1}{\sqrt{2}}, 0, \dots, 0\right),$$
$$v = \left(\frac{1}{\sqrt{2}}, -\frac{1}{\sqrt{2}}, 0, \dots, 0\right).$$

In the Figure 5.1 we show two elements of the Haar basis of length 16.



Figure 5.1: $R_8 u$ and $R_8 v$

2. Shannon first-stage wavelet basis.

Suppose N divisible by 4. We define,

$$\hat{u}(n) = \begin{cases} \sqrt{2} & n = 0, \dots, \frac{N}{4} - 1, \frac{3N}{4}, \dots, N - 1\\ 0 & n = \frac{N}{4}, \dots, \frac{3N}{4} - 1 \\ \end{pmatrix}$$
$$\hat{v}(n) = \begin{cases} 0 & n = 0, \dots, \frac{N}{4} - 1, \frac{3N}{4}, \dots, N - 1\\ \sqrt{2} & n = \frac{N}{4}, \dots, \frac{3N}{4} - 1 \\ \end{cases}$$

We observe that A(n) are unitary.

$$A\left(\frac{N}{4}\right) = \frac{1}{\sqrt{2}} \begin{pmatrix} \sqrt{2} & 0 & \cdots & 0\\ 0 & \cdots & 0 & \sqrt{2} \end{pmatrix}$$

3. Real Shannon first-stage wavelet basis.

Suppose N divisible by 4. We define,

$$\hat{u}(n) = \begin{cases} \sqrt{2} & n = 0, \dots, \frac{N}{4} - 1, \frac{3N}{4} + 1, \dots, N - 1 \\ i & n = \frac{N}{4} \\ -i & n = \frac{3N}{4} \\ 0 & otherwise \\ . \end{cases}$$

$$\hat{v}(n) = \begin{cases} 0 & n = 0, \dots, \frac{N}{4} - 1, \frac{3N}{4} + 1, \dots, N - 1 \\ 1 & n = \frac{N}{4} \\ 1 & n = \frac{3N}{4} \\ \sqrt{2} & otherwise \\ . \end{cases}$$

We observe that A(n) are unitary.

$$A\left(\frac{N}{4}\right) = \frac{1}{\sqrt{2}} \begin{pmatrix} i & 1\\ -i & 1 \end{pmatrix}$$

In the Figure 5.2 we show two elements of the Shannon real basis of length 64.



Figure 5.2: $R_{32}u$ and $R_{32}v$

4. Daub6 wavelet basis.

Ingrid Daubechies constructed families of wavelets that are very well localized in space rather than in frequency. We assume that N is divisible by 2^p , for some positive integer p, and that $N/2^p > 6$. Let M = N/2. Our goal is to construct $u \in \ell^2(\mathbb{Z}_N)$ such that u has only six nonzero components and satisfies the first equation in Theorem 5.11 and then we apply Lemma 5.13 to find v.

We begin with the trivial identity

$$\left[\cos^2\left(\frac{\pi n}{N}\right) + \sin^2\left(\frac{\pi n}{N}\right)\right]^5 = 1 \text{ for all } n.$$

Expanding this out, we have

$$\cos^{10}\left(\frac{\pi n}{N}\right) + 5\cos^8\left(\frac{\pi n}{N}\right)\sin^2\left(\frac{\pi n}{N}\right) + 10\cos^6\left(\frac{\pi n}{N}\right)\sin^4\left(\frac{\pi n}{N}\right)$$

$$(5.1) + 10\cos^4\left(\frac{\pi n}{N}\right)\sin^6\left(\frac{\pi n}{N}\right) + 5\cos^2\left(\frac{\pi n}{N}\right)\sin^8\left(\frac{\pi n}{N}\right) + \sin^{10}\left(\frac{\pi n}{N}\right) = 1.$$

Define

$$b(n) = \cos^{10}\left(\frac{\pi n}{N}\right) + 5\cos^8\left(\frac{\pi n}{N}\right)\sin^2\left(\frac{\pi n}{N}\right) + 10\cos^6\left(\frac{\pi n}{N}\right)\sin^4\left(\frac{\pi n}{N}\right).$$

Note that

$$\cos\left(\frac{\pi(n+M)}{N}\right) = \cos\left(\frac{\pi n}{N} + \frac{\pi}{2}\right) = -\sin\left(\frac{\pi n}{N}\right)$$

and similarly

$$\sin\left(\frac{\pi(n+M)}{N}\right) = \cos\left(\frac{\pi n}{N}\right).$$

Hence

$$b(n+M) = 10\cos^4\left(\frac{\pi n}{N}\right)\sin^6\left(\frac{\pi n}{N}\right) + 5\cos^2\left(\frac{\pi n}{N}\right)\sin^8\left(\frac{\pi n}{N}\right) + \sin^{10}\left(\frac{\pi n}{N}\right).$$

Thus by equation 5.1

$$b(n) + b(n+M) = 1$$
 for all n .

We select $u \in \ell^2(\mathbb{Z}_N)$ so that

(5.2)
$$|\hat{u}(n)|^2 = 2b(n).$$

Then we have

$$|\hat{u}(n)|^2 + |\hat{u}(n+M)|^2 = 2$$
 for all $n = 0, \dots, M-1$

We could obtain equation 5.2 by setting $u(n) = (\sqrt{2b(n)})$. But here we do not have a vector u with only six nonzero values. We write

$$b(n) = \cos^{6}\left(\frac{\pi n}{N}\right) \left[\cos^{4}\left(\frac{\pi n}{N}\right) + 5\cos^{2}\left(\frac{\pi n}{N}\right)\sin^{2}\left(\frac{\pi n}{N}\right) + 10\sin^{4}\left(\frac{\pi n}{N}\right)\right]$$
$$= \cos^{6}\left(\frac{\pi n}{N}\right) \left[\left(\cos^{2}\left(\frac{\pi n}{N}\right) - \sqrt{10}\sin^{2}\left(\frac{\pi n}{N}\right)\right)^{2} + \left(5 + 2\sqrt{10}\right)\cos^{2}\left(\frac{\pi n}{N}\right)\sin^{2}\left(\frac{\pi n}{N}\right)\right].$$

Define $\hat{u} \in \ell^2(\mathbb{Z}_N)$ by

$$\hat{u}(n) = \sqrt{2}e^{-2\pi i n/N} \cos^3\left(\frac{\pi n}{N}\right) \left[\cos^2\left(\frac{\pi n}{N}\right) - \sqrt{10} \sin^2\left(\frac{\pi n}{N}\right) + i\sqrt{5 + 2\sqrt{10}} \cos\left(\frac{\pi n}{N}\right) \sin\left(\frac{\pi n}{N}\right) \right].$$

By applying Euler's formula and the double angle identities, we have

$$\hat{u}(n) = \sqrt{2}e^{-2\pi i 4/N} e^{3\pi i n/N} \left(\frac{e^{i\pi n/N} + e^{-i\pi n/N}}{2}\right)^3 \\ \left[\frac{1}{2}\left(1 + \cos\left(\frac{2\pi n}{N}\right)\right) - \frac{\sqrt{10}}{2}\left(1 - \cos\left(\frac{2\pi n}{N}\right)\right) + i\frac{\sqrt{5 + 2\sqrt{10}}}{2}\sin\left(\frac{2\pi n}{N}\right)\right].$$

To simplify, $a = 1 - \sqrt{10}$, $b = 1 + \sqrt{10}$ and $c = \sqrt{5 + 2\sqrt{10}}$.

$$\hat{u}(n) = \frac{\sqrt{2}}{8} e^{-2\pi i 4n/N} \left(e^{2\pi i n/N} + 1 \right)^3 \left[\frac{a}{2} + \frac{b}{4} \left(e^{2\pi i n/N} + e^{-2\pi i n/N} \right) + \frac{c}{4} \left(e^{2\pi i n/N} - e^{-2\pi i n/N} \right) \right].$$

At this point we can see that

$$\hat{u}(n) = \sum_{k=0}^{5} u(k) e^{-2\pi i k n/N},$$

for some numbers u(0), u(1), u(2), u(3), u(4), u(5). If we calculate them, we obtain

$$u = (u(0), u(1), u(2), u(3), u(4), u(5), 0, \dots, 0)$$

= $\frac{\sqrt{2}}{32} (b + c, 2a + 3b - 3c, 6a + 4b + 2c, 6a + 4b - 2c, 2a + 3b - 3c, b - c, 0, \dots, 0).$

Finally, with Lemma 5.13 we obtain the vector v

$$v = (-u(1), u(0), 0, \dots, 0, -u(5), u(4), -u(3), u(2))$$

In the Figure 5.3 we show two elements of the Daubechies basis of length 64.



Figure 5.3: $R_{32}u$ and $R_{32}v$

We can compare Figure 5.2 with Figure 5.3 with same length.

Lemma 5.13. Suppose $M \in \mathbb{N}$, N = 2M and $u \in \ell^2(\mathbb{Z}_N)$ is such that $\{R_{2k}u\}_{k=0}^{M-1}$ is an orthonormal set with M elements. Define $v \in \ell^2(\mathbb{Z}_N)$ by

$$v(k) = (-1)^{k-1} \overline{u(1-k)} \ \forall k.$$

Then $\{R_{2k}u\}_{k=0}^{M-1} \bigcup \{R_{2k}v\}_{k=0}^{M-1}$ is a first-stage wavelets basis for $\ell^2(\mathbb{Z}_N)$.

Proof. We have to prove that the second and third conditions of Theorem 5.11 are satisfied.

First we will define $\hat{v}(m)$ and $\hat{v}(m+M)$:

1.

$$\begin{split} \hat{v}(m) &= \sum_{n=0}^{N-1} v(n) e^{-2\pi i m n/N} = \sum_{n=0}^{N-1} (-1)^{n-1} \overline{u(1-n)} e^{-2\pi i m n/N} \\ &\stackrel{k=1-n}{\longrightarrow} \sum_{n=0}^{N-1} (-1)^{-k} \overline{u(k)} e^{-2\pi i m (1-k)/N} \\ &= e^{-2\pi i m/N} \sum_{k=0}^{N-1} \overline{u(k)} (e^{-i\pi})^{-k} e^{2\pi i m k/N} = e^{-2\pi i m/N} \sum_{k=0}^{N-1} \overline{u(k)} e^{i\pi k + 2\pi i m k/N} \\ &= e^{-2\pi i m/N} \sum_{k=0}^{N-1} u(k) e^{i\pi k + 2\pi i m k/N} = e^{-2\pi i m/N} \overline{u(m+M)}. \end{split}$$

2.

$$\hat{v}(m+M) = e^{-2\pi i (m+M)/N} \overline{\hat{u}(m+2M)} = e^{-2\pi i m/N} e^{-2\pi i M/N} \hat{u}(m+N)$$
$$= -e^{-2\pi i m/N} \hat{u}(m).$$

Now we can prove the conditions

1.
$$|\hat{v}(m)|^2 + |\hat{v}(m+M)|^2 = |\hat{v}(m+M)|^2 + |\hat{v}(m)|^2 = 2.$$

2.

$$\hat{u}(m)\overline{\hat{v}(m)} + \hat{u}(m+M)\overline{\hat{v}(m+M)} = \hat{u}(m)e^{2\pi i m/N}\hat{u}(m+M) + \hat{u}(m+M)(-1)e^{2\pi i m/N}\hat{u}(m) = 0.$$

Suppose that $B = \{R_{2k}u\}_{k=0}^{M-1} \bigcup \{R_{2k}v\}_{k=0}^{M-1}$ is a first-stage wavelet basis. The change of basis of B to the euclidean basis E, is the matrix U, with columns the

vectors $v, R_2 v, \ldots, R_{N-2} v, u, R_2 u, \ldots, R_{N-2} u$. *B* is orthonormal, then *U* is a unitary matrix, so the change of basis from *E* to *B* is the matrix $U^{-1} = U^*$. Thus,

$$[z]_B = \begin{bmatrix} z * \tilde{v}(0) \\ \vdots \\ z * \tilde{v}(N-2) \\ z * \tilde{u}(0) \\ \vdots \\ z * \tilde{u}(N-2) \end{bmatrix}$$

where $[z]_B$ is the vector $z \in \ell^2(\mathbb{Z}_N)$ in B basis.

Definition 5.14. Suppose $M \in \mathbb{N}$, N = 2M and $u \in \ell^2(\mathbb{Z}_N)$. We define the Downsampling operator $D : \ell^2(\mathbb{Z}_N) \longrightarrow \ell^2(\mathbb{Z}_M)$

$$D(z)(n) = z(2n), \text{ for } n = 0, \dots, M-1$$

In other words, if z = (z(0), z(1), ..., z(N-1)) the downsampling operator is often denoted $\downarrow 2$ in diagrams. The calculation of $[z]_B$ is represented in Figure 5.4.

We have seen that we can compute the E to B change of basis quickly. Now we are going to compute the B to E change of basis. This can be obtained by multiplying by the matrix U, but it takes $N \times N$ multiplications, and this is slow. In Figure 5.4 we see a fast procedure basek on the *filter bank* approach.

Definition 5.15. A filter bank is an array of band-pass filters that separates the input signal into multiple components, each one carrying a simple frequency subband of the original signal.

$$z \xrightarrow{\qquad \quad \\ \ast \tilde{v} \xrightarrow{\qquad \quad \\ } z \ast \tilde{v} \xrightarrow{\qquad \quad \\ } \downarrow 2 \xrightarrow{\qquad \quad \\ } D(z \ast \tilde{v}) = \begin{bmatrix} z \ast \tilde{v}(0) \\ \vdots \\ z \ast \tilde{v}(N-2) \\ \vdots \\ z \ast \tilde{u}(0) \\ \vdots \\ z \ast \tilde{u}(N-2) \end{bmatrix} \approx \begin{bmatrix} z \ast \tilde{v}(0) \\ \vdots \\ z \ast \tilde{u}(0) \\ \vdots \\ z \ast \tilde{u}(N-2) \end{bmatrix} = [z]_B$$

Figure 5.4

Definition 5.16. Suppose $M \in \mathbb{N}$ and N = 2M. Define $U : \ell^2(\mathbb{Z}_M) \longrightarrow \ell^2(\mathbb{Z}_N)$ for $z \in \ell^2(\mathbb{Z}_M)$,

$$U(z)(n) = \begin{cases} z(n/2) & \text{if } n \text{ is even} \\ 0 & \text{if } n \text{ is odd} \end{cases}$$

The operator U is called the upsampling operator. It is denoted by $\uparrow 2$.

The *upsampling operator* doubles the size of a vector by inserting a 0 between any two adjacent values.

Remark 5.17.

1.

$$(5.3) D \circ U(z) = z,$$

2.

(

5.4)
$$U \circ D(z) = \frac{1}{2}(z + z^*).$$

To regain z from the output of the left filter bank in Figure 5.4, we follow up with a right filter bank as in the right half of Figure 5.5 Here $s, t \in \ell^2(\mathbb{Z}_N)$ are unknown. The output of the upper branch of Figure 5.5 is $\tilde{t} * U(D(z * \tilde{v}))$ and the output of the lower branch is $\tilde{s} * U(D(z * \tilde{u}))$. The lemma below gives conditions under which the sum of these outputs is always the original input z. When this happens we say we have *perfect reconstruction* in the filter bank. Note that we do not necessarily assume the conditions of Theorem 5.11.



Figure 5.5

Lemma 5.18. Suppose $M \in \mathbb{N}$, N = 2M and $u, v, s, t \in \ell^2(\mathbb{Z}_N)$. For $n = 0, \ldots, N - 1$, let A(n) be the system matrix (def.) for u and v. Then we have perfect reconstruction in Figure 5.5, that is,

$$\tilde{t} * U(D(z * \tilde{v})) + \tilde{s} * U(D(z * \tilde{u})) = z$$

for all $z \in \ell^2(\mathbb{Z}_N)$, if and only if

$$A(n) \cdot \begin{bmatrix} s(n) \\ t(n) \end{bmatrix} = \begin{bmatrix} \sqrt{2} \\ 0 \end{bmatrix}$$

for all n = 0, ..., N - 1. In the case that A(n) is unitary, this simplifies to $\hat{t}(n) = \overline{\hat{v}(n)}$ and $\hat{s}(n) = \overline{\hat{u}(n)}$. If A(n) is unitary for all n, then $t = \tilde{v}$ and $s = \tilde{u}$.

Proof. By properties of the upsampling and downsampling operator we have that

$$U(D(z * \tilde{v})) = \frac{1}{2} \left((z * \tilde{v}) + (z * \tilde{v})^* \right)$$

and similarly with v replaced by u.

Remark 5.19. $z^*(n) = (-1)^n z(n) \ \forall n \Rightarrow (z^*)\hat{}(n) = \hat{z}(n+M) \ \forall n.$

Taking Fourier transform,

$$U(D(z * \tilde{v}))(n) = \frac{1}{2}(\hat{z}(n)\overline{\hat{v}(n)} + (z * \tilde{v})(n + M))$$
$$= \frac{1}{2}(\hat{z}(n)\overline{\hat{v}(n)} + (\hat{z}(n + M)\overline{\hat{v}(n + M)}) \forall n$$

and similarly with v replaced by u.

Let us see that

$$\tilde{t} * U(D(z * \tilde{v})) + \tilde{s} * U(D(z * \tilde{u})) = z, \ \forall z \in \ell^2(\mathbb{Z}_N).$$

We will prove that

$$\left(\tilde{t} * U(D(z * \tilde{v})) + \tilde{s} * U(D(z * \tilde{u}))\right)(n) = \hat{z}(n), \ \forall z \in \ell^2(\mathbb{Z}_N).$$

$$\begin{split} &\left[\tilde{t}*U(D(z*\tilde{v}))+\tilde{s}*U(D(z*\tilde{u}))\right](n)\\ &=\overline{\hat{t}(n)}\frac{1}{2}\left[\hat{z}(n)\hat{v}(n)+\hat{z}(n+M)\hat{v}(n+M)\right]+\overline{\hat{s}(n)}\frac{1}{2}\left[\hat{z}(n)\hat{u}(n)+\hat{z}(n+M)\hat{u}(n+M)\right]\\ &=\hat{z}(n)\left[\frac{1}{2}\left(\overline{\hat{t}(n)}\overline{\hat{v}(n)}+\overline{\hat{s}(n)}\hat{u}(n)\right)\right]+\frac{\hat{z}(n+M)}{2}\left[\overline{\hat{t}(n)}\overline{\hat{v}(n+M)}+\overline{\hat{s}(n)}\hat{u}(n+M)\right]\\ &=\hat{z}(n)\ \forall n. \end{split}$$

This happens if and only if,

$$\hat{t}(n)\hat{v}(n) + \hat{s}(n)\hat{u}(n) = 2$$

and

$$\hat{t}(n)\hat{v}(n+M) + \hat{s}(n)\hat{u}(n+M) = 0,$$

and it is equivalent to

$$\begin{pmatrix} \hat{u}(n) & \hat{v}(n) \\ \hat{u}(n+M) & \hat{v}(n+M) \end{pmatrix} \begin{pmatrix} \hat{s}(n) \\ \hat{t}(n) \end{pmatrix} = \begin{pmatrix} 2 \\ 0 \end{pmatrix}$$

Notice that

$$\sqrt{2}A(n) = \begin{pmatrix} \hat{u}(n) & \hat{v}(n) \\ \hat{u}(n+M) & \hat{v}(n+M) \end{pmatrix}$$

In the case that A(n) is unitary, A(n) is invertible and $(A(n))^{-1} = A^*(n)$, so solving

$$\begin{pmatrix} \hat{s}(n)\\ \hat{t}(n) \end{pmatrix} = A^*(n) \begin{pmatrix} \sqrt{2}\\ 0 \end{pmatrix}$$

From this equality we obtain $\hat{s}(n) = \overline{\hat{u}(n)}$ and $\hat{t}(n) = \overline{\hat{v}(n)}$. Then, if A(n) is unitary for all $n, s(n) = \tilde{u}$ and $t(n) = \tilde{v}$. Notice that $\tilde{s} = u, \tilde{t} = v$.

Remark 5.20. If u, v are wavelets, then



Figure 5.6

If $N = \dot{4}$, we will do 2 steps, if $N = \dot{8}$, we will do 3 steps and so on.



Figure 5.7

Definition 5.21. If $N = 2^p$, we define a sequence of wavelet filters of stage p with a sequence of vectors $(u_1, v_1), \ldots, (u_p, v_p)$ such that

1. $(u_l, v_l) \in \ell^2(\mathbb{Z}_{N/2^{l-1}}), l = 1, \dots, p.$

2.

$$A_l(n) = \begin{pmatrix} \hat{u}_l(n) & \hat{v}_l(n) \\ \hat{u}_l(n + \frac{N}{2^l}) & \hat{v}_l(n + \frac{N}{2^l}) \end{pmatrix}$$

for $n = 0, \ldots, \frac{N}{2^l} - 1$ the matrix are unitary.

5.3 Matlab Implementation

5.3.1 Haar Wavelets

```
_____ Code _____
function H=Haar1D(x)
%Supose rows=1 (x is a vector)
[rows,cols] = size(x);
% Check if size = pow(2), otherwise call padding function
[new_size,N] = padding(cols);
x_new = zeros(rows,new_size);
x_new(1:rows,1:cols)=x;
% Calcul Haar transform
v = x_new';
A = CalculTransf(pow2(N-1));
v_new = A*v;
H=[v_new(1:pow2(N-1)), v_new(pow2(N-1)+1:end)];
%%%%%% Auxiliar function %%%%%%%%%%%%
%%% Matrix transformation
function A=CalculTransf(dim)
for i=1:dim
   S(i,i*2-1:i*2)=1/sqrt(2);
```

end

```
for i=1:dim
    D(i,i*2-1)=1/sqrt(2);
    D(i,i*2)=-1/sqrt(2);
end
A = [S;D];
%%% Padding
function [m,j]=padding(n)
m=2;
j=1;
while n>m
    m=m*2;
    j=j+1;
```

end



Figure 5.8: U matrix of Haar basis

5.3.2 Shannon Wavelets

```
____ Code _____
function z2=Shannon(z)
[rows,cols] = size(z);
%Check if m is multiple of 4
[m]=multiple4(cols);
v = zeros(1, m-cols);
z = [z v];
U = basisShannon(m);
for i=1:m
    z2(i) = sum(U(:,i)'.*z);
end
%% Auxiliar function
function U=basisShannon(m)
%% Calcule of Shannon basis in dimension m
fu = zeros(1,m);
fv = zeros(1,m);
fu(1:m/4) = sqrt(2);
fu(3*m/4 + 2:end) = sqrt(2);
fu(m/4 + 1)=j;
fu(3*m/4 + 1)=-j;
fv(m/4+2:3*m/4) = sqrt(2);
fv(3*m/4 + 1)=1;
fv(m/4 + 1)=1;
u=ifft(fu);
```

```
v=ifft(fv);
%% Calcule matrix U
for i=1:m/2
    U(:,i) = circshift(u,[1 (i-1)*2])';
end
for i=1:m/2
    U(:,i+m/2) = circshift(v,[1 (i-1)*2])';
end
```



Figure 5.9: U matrix of Shannon basis, real part



Figure 5.10: U matrix of Shannon basis, imaginary part
5.3.3 Daubechies Wavelets

We will compute the *Daubechies* 4 and *Daubechies* 6.

```
_Code _
function Daub = Daub4(x)
%% Define alpha & beta
a1 = (1+sqrt(3))/(4*sqrt(2));a2 = (3+sqrt(3))/(4*sqrt(2));
a3 = (3-sqrt(3))/(4*sqrt(2));a4 = (1-sqrt(3))/(4*sqrt(2));
b1 = a4; b2 = -a3; b3 = a2; b4 = -a1;
[rows,cols]=size(x);
% Check if size = pow(2), otherwise call padding function
[new_size,N] = padding(cols);
x_new = zeros(rows,new_size);
x_new(1:rows,1:cols)=x;
a = zeros(1,new_size);
a(1:4) = [a1 \ a2 \ a3 \ a4];
b = zeros(1,new_size);
b(1:4)=[b1 b2 b3 b4];
% Calcule of matrix
m=0;
for i=1:new_size/2
    S(i,:)=circshift(a,[1 m]);
    m=m+1;
end
m=0;
for i=1:new_size/2
    D(i,:)=circshift(b,[1 m]);
```

```
m=m+1;
end
A=[S;D];
Daub=A*x_new';
%%% Padding
function [m,j]=padding(n)
m=2;
j=1;
while n>m
m=m*2;
j=j+1;
end
```



Figure 5.11: U matrix Daubechies4 basis

Daubechies 6 function.

Code _

```
function U = Daub6(rows,cols)
a=1-sqrt(10);
b=1+sqrt(10);
c=sqrt(5+2*sqrt(10));
u = zeros(1,cols);
v = zeros(1,cols);
u(1) = b+c;
u(2) = 2*a + 3*b + 3*c;
u(3) = 6*a + 4*b + 2*c;
u(4) = 6*a + 4*b - 2*c;
u(5) = 2*a + 3*b - 3*c;
u(6) = b-c;
u = u.*(sqrt(2)/32);
v(1) = -u(2);
v(2) = u(1);
v(end)=u(3);
v(end-1)=-u(4);
v(end-2)=u(5);
v(end-3)=-u(6);
%% Calcule matrix U
for i=1:cols/2
    U(:,i) = circshift(u,[1 2*(i-1)]).';
    U(:,i+cols/2) = circshift(v,[1 2*(i-1)]).';
end
```



Figure 5.12:U matrix of Daubechies6 basis

Chapter 6

Applications: Compression and De-noising

In this chapter we will deal with wavelets basis for compression and de-noising signal. The reason why wavelets basis are used for these applications is because they are well localized in frequency, and moreover because Daubechies wavelets are also well localized in space (as we will see in the examples of image compression). For a better comprehension of one dimensional signal processing, to analyse the spectrum of music, see [5].

For each $N \in \mathbb{N}$ one can use *ibechies's DN wavelets* for $\ell^2(\mathbb{Z})$, for which u has N nonzero components, to construct supported wavelets for $L^2(\mathbb{R})$. As N increases, *Daubechies's DN wavelets* have larger support intervals. However, this is compensated in many applications by the facts that DN tend to have more cancellation and more smoothness.

The smoothness of *bechies's wavelets* grows linearly with their support, and linear growth is the best. For both *Daubechies wavelets transforms* and the corresponding scaling function require just a few values from the signal to get a good approximation, and that is *Daubechies* provide us with a powerful tool for signal processing.

6.1 Signal Compression

Wavelet transform produces a large number of values having zero, or near zero, magnitudes. The compression signal is obtained by applying a wavelet transform, and getting only the largest values obtained. After that, we apply the inverse transform and we will get the compression signal. Here, we will not deal with the encoding of those pixels that are set to zero.

To compare the results with the different wavelets basis we introduce the definition of the error that will measure the output signals.

Definition 6.1. Mean Square Error(MSE) is defined as

$$MSE = \frac{1}{N} \sum_{i,j} (f(i,j) - g(i,j))^2,$$

where N is the total number of elements, and the sum over i, j is the sum over all the elements in the signal, f is the original signal and g is its compression.

In the examples the relative error is displayed, and it is calculated by dividing the MSE with the norm of the signal

$$Error = N \frac{MSE}{norm(f)},$$

where f is the given signal.

Example 6.2. We introduce two examples of 1 dimensional signal compression.

1.

$$z(n) = \begin{cases} 0 & 0 \le n \le 127\\ \sin\left(\frac{|n-128|^{1.7}}{128}\right) & 128 \le n \le 255\\ 0 & 256 \le n \le 383\\ \sin\left(\frac{|n-128|^2}{128}\right) & 384 \le n \le 447\\ 0 & 448 \le n \le 511 \end{cases}$$



Figure 6.1

2.

$$z(n) = \begin{cases} 1 - \frac{n}{64} & 0 \le n \le 63\\ 0 & 64 \le n \le 255\\ 5 - \frac{n}{64} & 256 \le n \le 319\\ 0 & 320 \le n \le 447 \end{cases}$$



Figure 6.2



Figure 6.3

In this example we can see that in Figure 6.2 it seems that the Fourier transform compress better the signal, but as we can see in Figure 6.3, where we have taken more coefficients in the transformation, the Fourier compression is still having problems in those parts where have large changes, and not like wavelets transforms. With only the 20% of the coefficients the Daubechies6 has an error of 10^{-16} .

Example 6.3. Here, we have the image of barbara.png. It has been compressed by taking the 10% of the largest coefficients of each wavelet and Fourier transform. If we compare the results by using Fourier and wavelets, we see that the outputs in Haar and Daubechies wavelets are better than the first one, as explained in Chapter 5.



(a) Original image.



(c) Haar 10% coeff.



(b) Fourier 10% coeff.



(d) Daubechies 610% coeff.

Figure 6.4: Compression *barbara.png* image.



Figure 6.5

6.1.1 Matlab implementation

This function calculates Fourier, Daubechies6 and Shannon wavelet transformation.

```
_ Code _
function compressWavelets(x,K)
%%
%%
  Compress signal x taking the higher
%%
  values in diferent transform
%%
  wavelets
%%
[rows,cols]=size(x);
```

 $v_{max} = [max(x)];$ $v_{\min} = [\min(x)];$

```
% Fast Fourier
F = fft(x);
F_err = F;
[order,v_f]=sort(abs(F),'descend');
F(v_f(K+1:end))=0;
Fourier = ifft(F);
v_max = [v_max max(real(Fourier))];
v_min = [v_min min(real(Fourier))];
error_Fourier = norm(real(Fourier) - x)/norm(x)
```

```
% D6
U = TransfD6(rows,cols);
D = D6(x,U);
D_err = D;
```

```
[order_d,v_d] = sort(abs(D),'descend');
D(v_d(K+1:end))=0;
Daub6=ID6(D,U);
v_max = [v_max max(real(Daub6))];
v_min = [v_min min(real(Daub6))];
error_Daub6 = norm(-(Daub6) + x)/norm(x)
```

```
% Shannon transform
Us = TransfSh(rows,cols);
S = Shannon2(x,Us);
S_err = S;
```

[order_s,v_s]=sort(abs(S),'descend');

```
S(v_s(K+1:end))=0;
Shannon=IShannon(S,Us);
v_max = [v_max max(real(Shannon))];
```

```
v_min = [v_min min(real(Shannon))];
error_Shannon = norm(real(Shannon) - x')/norm(x)
%% Relative error for k=0:300
error_f=[];
error_d=[];
error_s=[];
norm_x = norm(x);
for k=100:-1:1
    F_err(v_f(k:end))=0;
    Fourier_err = ifft(F_err);
    error_f = [error_f norm(real(Fourier_err) - x)/norm_x];
      D_{err}(v_d(k:end))=0;
      Daub6_err = ID6(D_err, U);
      error_d = [error_d norm(real(Daub6_err) - x)/norm_x];
     S_err(v_s(k:end))=0;
     Shannon_err = IShannon(S_err,Us);
     error_s = [error_s norm(real(Shannon_err) - x')/norm_x];
end
%% Plot the results
figure;
subplot(3,2,1);plot(x);axis([0 length(x) min(v_min) max(v_max)]);
title('Signal');
subplot(3,2,2);plot(real(Fourier));
axis([0 length(x) min(v_min) max(v_max)]);
title(['Fourier with ',int2str(K),' higher coeff.']);
subplot(3,2,3);plot(real(Daub6));
axis([0 length(x) min(v_min) max(v_max)]);
```

```
title(['Daub6 with ',int2str(K),' higher coeff.']);
subplot(3,2,4);plot(real(Shannon));
axis([0 length(x) min(v_min) max(v_max)]);
title(['Shannon with ',int2str(K),' higher coeff.']);
subplot(3,2,5);plot(error_f(end:-1:1),'b');hold on;
plot(error_d(end:-1:1),'g');plot(error_s(end:-1:1),'r');hold off;
title('Red: Shannon. Green: Daub6. Blue: Fourier.');
```

Compression image by *Daubechies6*.

```
Code
function [fcompres,coef]=compress_image_d6(signal,level,eps)
[r,c]=size(signal);
[r,c]=size(signal);
%Order
F_abs=abs(signal);
[order,x]=ordenar_matriz(F_abs);
%Be carefull, now order is a vector!!
Fcompres=signal;
Fcompres=reshape(Fcompres,1,[]);
Fcompres(x(eps+1:end))= 0;
Fcompres = reshape(Fcompres,[r c]);
```

fcompres = InvD6_rkl(Fcompres,level);

Compression image by *Haar*.

Code _____

function [fcompres,coef]=compress_image_haar(signal,level,eps)

[r,c]=size(signal);

%Order

```
F_abs=abs(signal);
[order,x]=ordenar_matriz(F_abs);
%Be carefull, now order is a vector!!
Fcompres=signal;
Fcompres=reshape(Fcompres,1,[]);
Fcompres(x(eps+1:end))= 0;
Fcompres = reshape(Fcompres,[r c]);
fcompres = Inv_rkl(Fcompres,level);
```

Auxiliar function order matrix.

function [max_v,x]=ordenar_matriz(A)

[r,c]=size(A); B = reshape(A,1,[]); [max_v,x] = sort(B,'descend');

6.2 Signal De-noising

In this section we will obtain signal de-noising by observing the wavelet transformation of the signal, and thresholding it to remove the frequencies of noise. We have applied this same reasoning to separate frequencies of a given signal that has distinctives high frequencies and low frequencies.

Example 6.4. In both Figures 6.6 and 6.7 we have used the Daubechies6 transform of level 4 and 5 respectively. The processed signal is a sinusodial with Gaussian noise. If we consider the transform signal, we will see that noise is between the

interval [-3 3]. Then, by thresholding the transformation with this interval, we will obtain the desire signal. Notice that the more level we apply in the transform, the best results are obtained.



Figure 6.6: Signal De-noising.



Figure 6.7: Signal De-noising.

Example 6.5. Given a signal with two different frequencies, we attempt to separate them by thresholding the wavelet transform, obtaining the high frequencies and the low frequencies, and finally applying the inverse transform to each one and getting the two different signals.



Figure 6.8: Separate frequencies *barreja.wav*.



Figure 6.9: Separate frequencies *murcielate.wav*.

In the implementation of the program, we only have selected an array of lenght 1024 in the signal, because the size of the original signals exceeds the memory of Matlab program. Even so, the frequencies are well separated in this region of the array.

Example 6.6. Here we deal with image de-noising. In Figure 6.10 we have the picture of lena.png with Gaussian noise. If we look to Figure 6.10, we see that the histogram of the image has form of a Gaussian function. Then, if we calculate the

deviation of the original image, when we make the wavelet transform, we can delete the noise by thresholding with the value of its deviation.



Figure 6.10: Histogram of *lena.png* with Gaussian noise.

In this example we have de-noised the image using two different values of threshold. In Figure 6.11 we have used the deviation of the original image, which is equal to 0.21. In Figure 6.12 we have used three times the deviation of the image. We observe that in first figure the noise is not completely deleted, while in the second figure there is not noise, but all the image has been blured.



Figure 6.11: De-noising *lena.png* with Gaussian noise. Threshold: 0.21.



(a) Original image. (b) Haar (c) Daubechies6

Figure 6.12: De-noise *lena.png* with Gaussian noise. Threshold: 3*0.21.

6.2.1 Matlab implementation

Daubechies6 de-noising signal.

Daubechies6 separating frequencies.

Image de-noising by *Daubechies6*.

_____ Code _____

function denoise=denoise_image_d6(signal,level)

Fdenoise=signal; dev = std2((signal)); denoise2 = Fdenoise.*(abs(Fdenoise)> 3*dev); denoise = InvD6_rkl(denoise2,level);

Image de-noising by *Haar*.

_____ Code _____

function denoise=denoise_image_haar(signal,level)

```
Fdenoise=signal;
dev = std2(abs(signal));
denoise2 = Fdenoise.*(abs(Fdenoise)> 3*dev);
denoise = Inv_rkl(denoise2,level);
```

Bibliography

F

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